## Quantum Data Management From Theory to Opportunities

### **Rihan Hai, Shih-Han Hung, Sebastian Feld**











### Where it all started

#### **Simulating Physics with Computers**

**Richard P. Feynman** 

Department of Physics, California Institute of Technology, Pasadena, California 91107

Received May 7, 1981

#### 1. INTRODUCTION

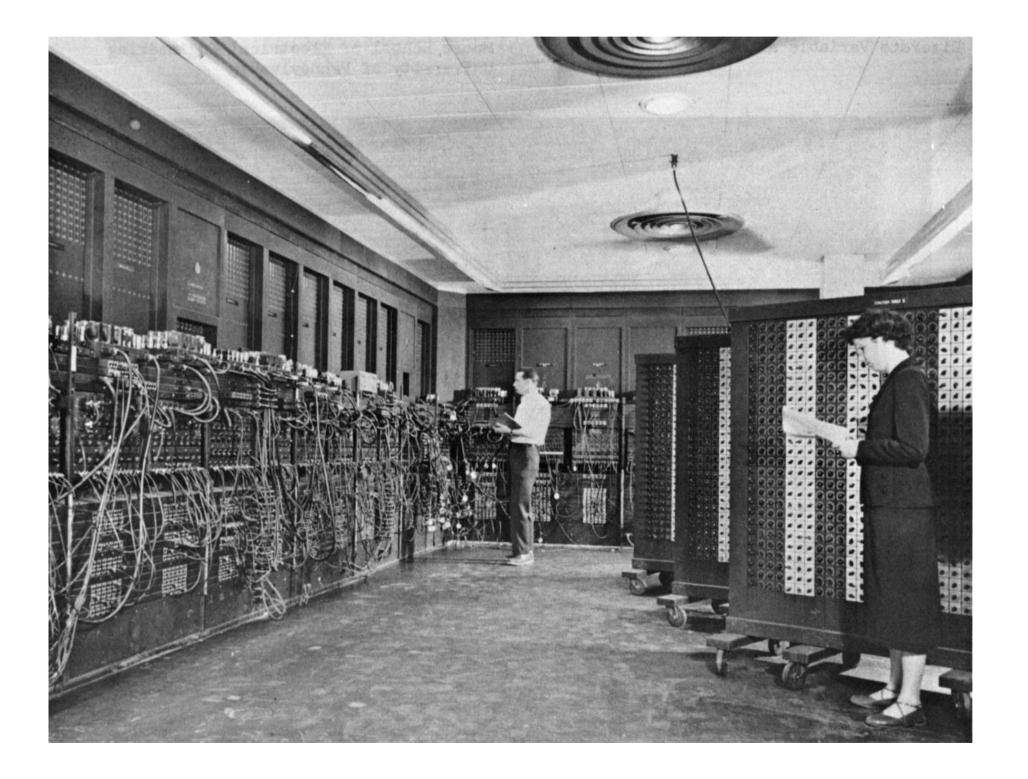
On the program it says this is a keynote speech—and I don't know what a keynote speech is. I do not intend in any way to suggest what should be in this meeting as a keynote of the subjects or anything like that. I have my own things to say and to talk about and there's no implication that anybody needs to talk about the same thing or anything like it. So what I want to talk about is what Mike Dertouzos suggested that nobody would talk about. I want to talk about the problem of simulating physics with computers and I mean that in a specific way which I am going to explain. The reason for doing this is something that I learned about from Ed Fredkin, and my entire interest in the subject has been inspired by him. It has to do with learning something about the possibilities of computers, and also something about possibilities in physics. If we suppose that we know all the physical laws perfectly, of course we don't have to pay any attention to computers. It's interesting anyway to entertain oneself with the idea that we've got something to learn about physical laws; and if I take a relaxed view here (after all I'm here and not at home) I'll admit that we don't understand everything.

The first question is, What kind of computer are we going to use to simulate physics? Computer theory has been developed to a point where it realizes that it doesn't make any difference; when you get to a *universal computer*, it doesn't matter how it's manufactured, how it's actually made. Therefore my question is, Can physics be simulated by a universal computer? I would like to have the elements of this computer *locally interconnected*, and therefore sort of think about cellular automata as an example (but I don't want to force it). But I do want something involved with the

be understood very well in analyzing the situation. And I'm not happy with all the analyses that go with just the classical theory, because nature isn't classical, dammit, and if you want to make a simulation of nature, you'd better make it quantum mechanical, and by golly it's a wonderful problem, because it doesn't look so easy. Thank you.



### **Back to the future**



en.wikipedia.org/wiki/ENIAC#/media/File:Glen\_Beck\_and\_Betty\_Snyder\_program\_the\_ENIAC\_in\_building\_328\_at\_the\_Ballistic\_Research\_Laboratory.jpg



## Lots of activities

Article | Published: 23 October 2019

#### Quantum supremacy using a programmable superconducting processor

Frank Arute, Kunal Arya, Ryan Babbush, Dave Bacon, Joseph C. Bardin, Rami Barends, Rupak Biswas, Sergio Boixo, Fernando G. S. L. Brandao, David A. Buell, Brian Burkett, Yu Chen, Zijun Chen, Ben Chiaro, Roberto Collins, William Courtney, Andrew Dunsworth, Edward Farhi, Brooks Foxen, Austin Fowler, Craig Gidney, Marissa Giustina, Rob Graff, Keith Guerin, ... John M. Martinis 🗠 🕇 Show authors

Nature 574, 505–510 (2019) Cite this article

#### Simulating key properties of lithium-ion batteries with a faulttolerant quantum computer

Alain Delgado, Pablo A. M. Casares, Roberto dos Reis, Moditaba Shokrian Zini, Roberto Campos, Norge Cruz-Hernández, Arne-Christian Voigt, Angus Lowe, Soran Jahangiri, M. A. Martin-Delgado, Jonathan E. Mueller, and Juan Miguel Arrazola

Phys. Rev. A 106, 032428 – Published 26 September 2022

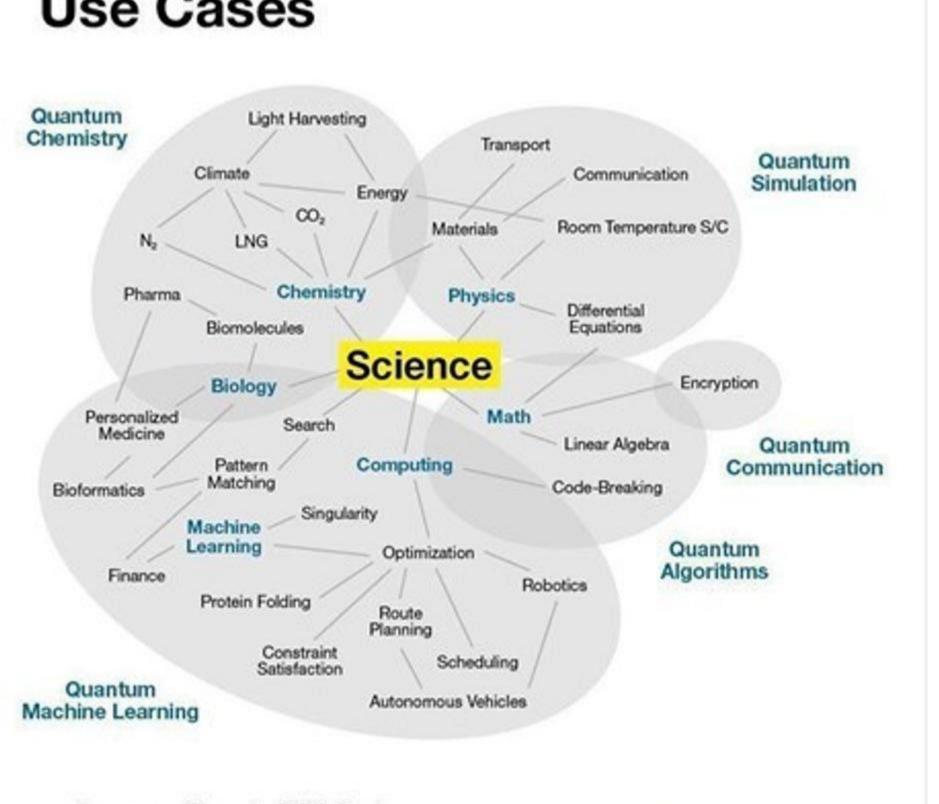
April 1, 2018

nature.com/articles/s41586-019-1666-5 journals.aps.org/pra/abstract/10.1103/PhysRevA.106.032428 writings.stephenwolfram.com/2018/04/buzzword-convergence-making-sense-of-guantum-neural-blockchain-ai/

#### **Buzzword Convergence: Making Sense of** Quantum Neural Blockchain AI

## Lots of (potential) use cases

#### Quantum Computing Use Cases



#### gartner.com/SmarterWithGartner

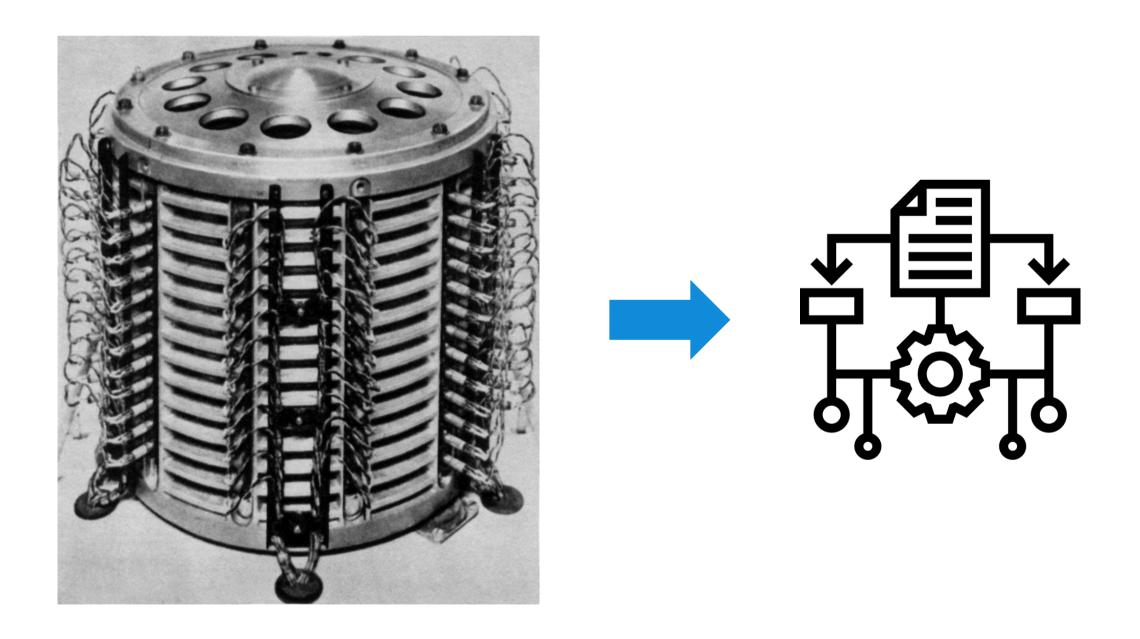
Source: Adapted from Pete Shadbolt and Jeremy O'Erlen © 2017 Gartner, Inc. and/or its attiliates. All rights reserved. Gartner is a registered trademark of Gartner, Inc. or its attiliates. FPL\_338248



forbes.com/sites/chuckbrooks/2021/03/21/the-emerging-paths-of-quantum-computing/?sh=66cac7936613

### Menu for today

# Entering the era of quantum computing, what is the future of data management?



#### Data management problem

Algorithm

de.wikipedia.org/wiki/Trommelspeicher#/media/Datei:Pamiec\_bebnowa\_1.jpg flaticon.com/de/kostenloses-icon/algorithmus\_1119005 forbes.com/sites/tiriasresearch/2023/11/28/quantum-computing-is-coming-faster-than-you-think/



#### Quantum computer

### **Contents and instructors**

Structure

- Introduction
- Fundamentals of quantum computing
- Data management using quantum computers
- Data management via quantum internet



**Rihan Hai** TU Delft, NL



Shih-Han Hung Academica Sinica, TW

#### [~5'] [~20'] [~35'] [~30']



Sebastian Feld TU Delft, NL

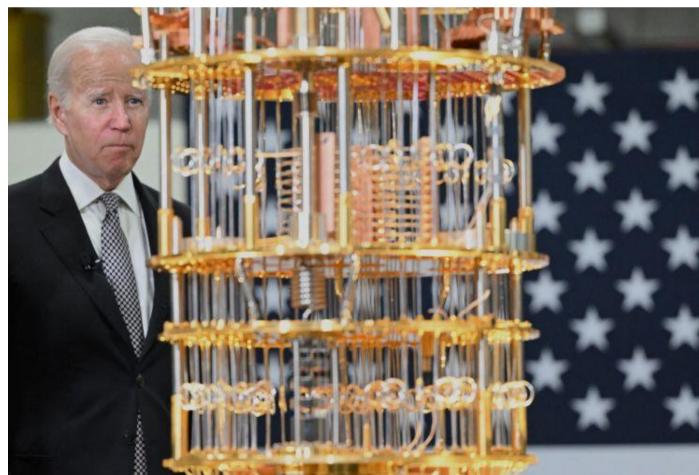
- Overview
- Quantum Gate Model
- Quantum Annealing
- Conclusion

- Overview
- Quantum Gate Model
- Quantum Annealing
- Conclusion

Golden chandelier plus VIP







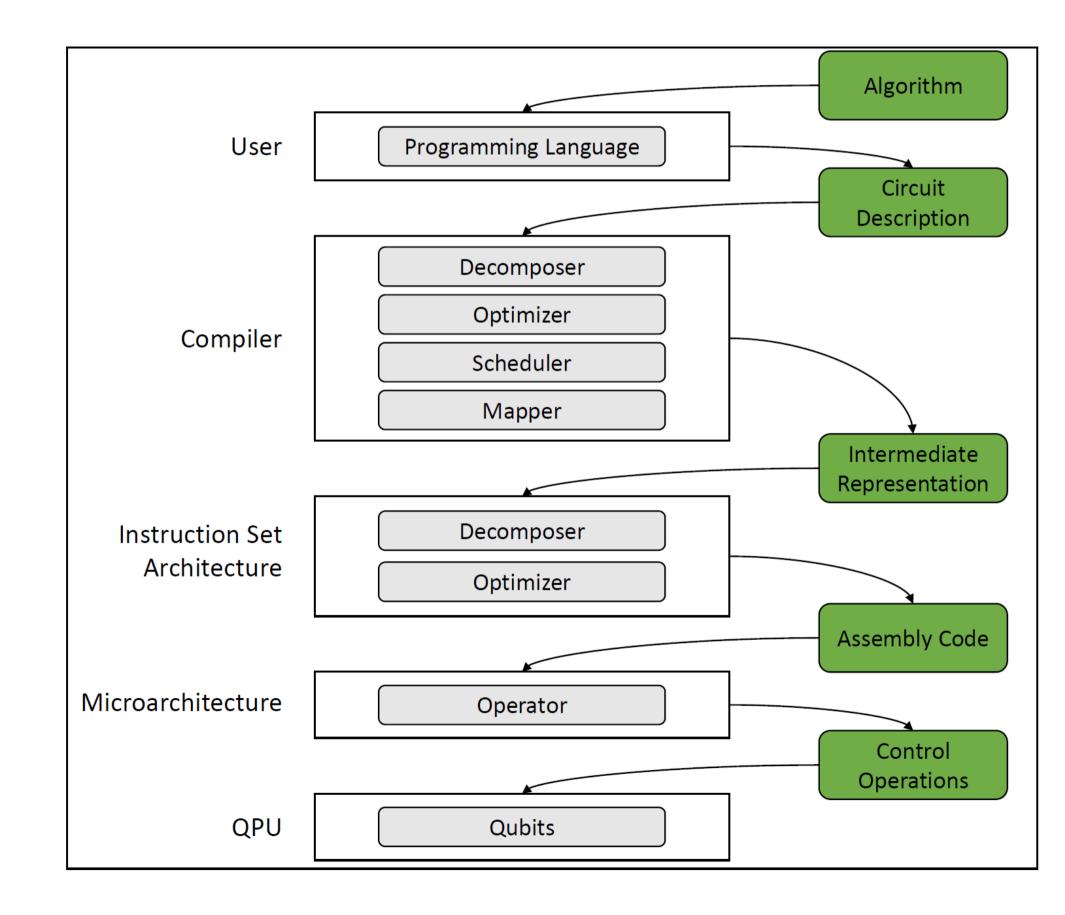
• Jokes aside: what is included and what is not?



#### Often hidden behind a service



Full-stack quantum computation

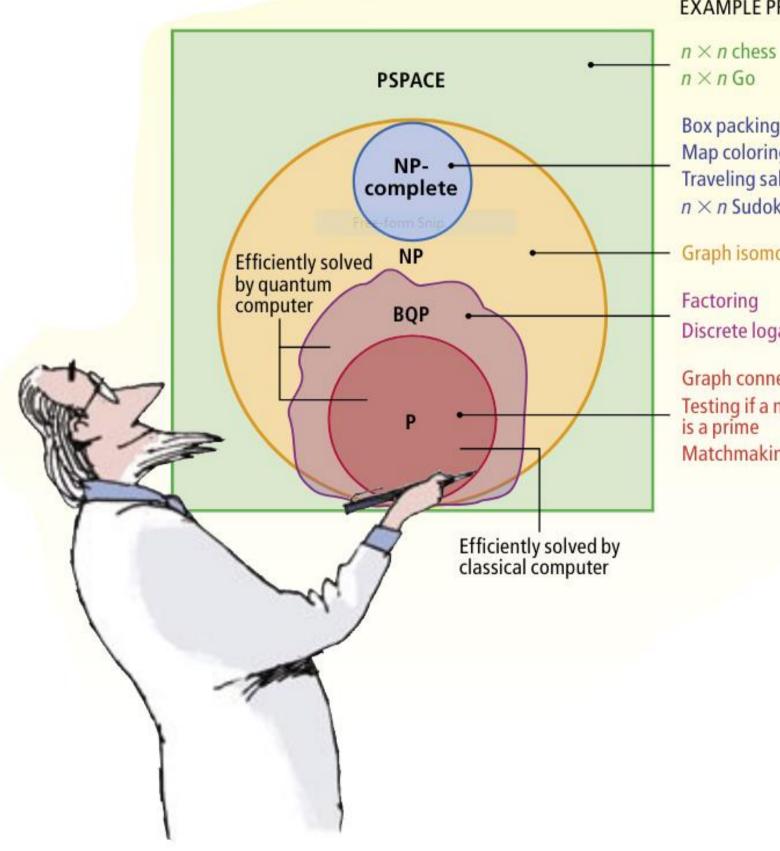


- It's not a magical wonder machine
- Quantum computing is turing-complete



## **Know your limits**

 Quantum computing is powerful, and we don't even know the boundaries



#### **EXAMPLE PROBLEMS** Box packing Map coloring Traveling salesman $n \times n$ Sudoku Graph isomorphism Harder Discrete logarithm Graph connectivity Testing if a number Matchmaking

## The holy grail

At least three things need to come together

No efficient classical solution

Valuable Problem

Efficient quantum solution

#### That's the goal!

## **Two different flavors**

Quantum Gate Model







Microsoft



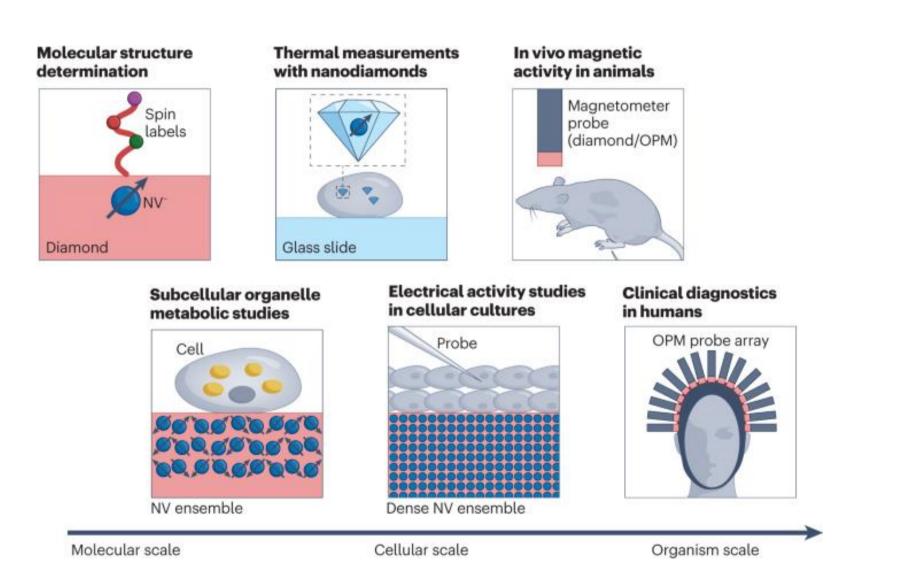


#### Quantum Annealing

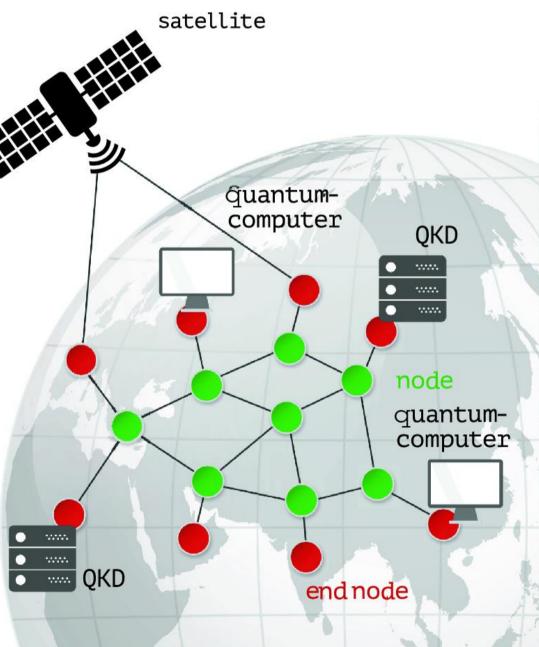
### D:VO/CThe Quantum Computing Company™

### And what else?

#### Quantum Sensing







#### Quantum Internet

Overview

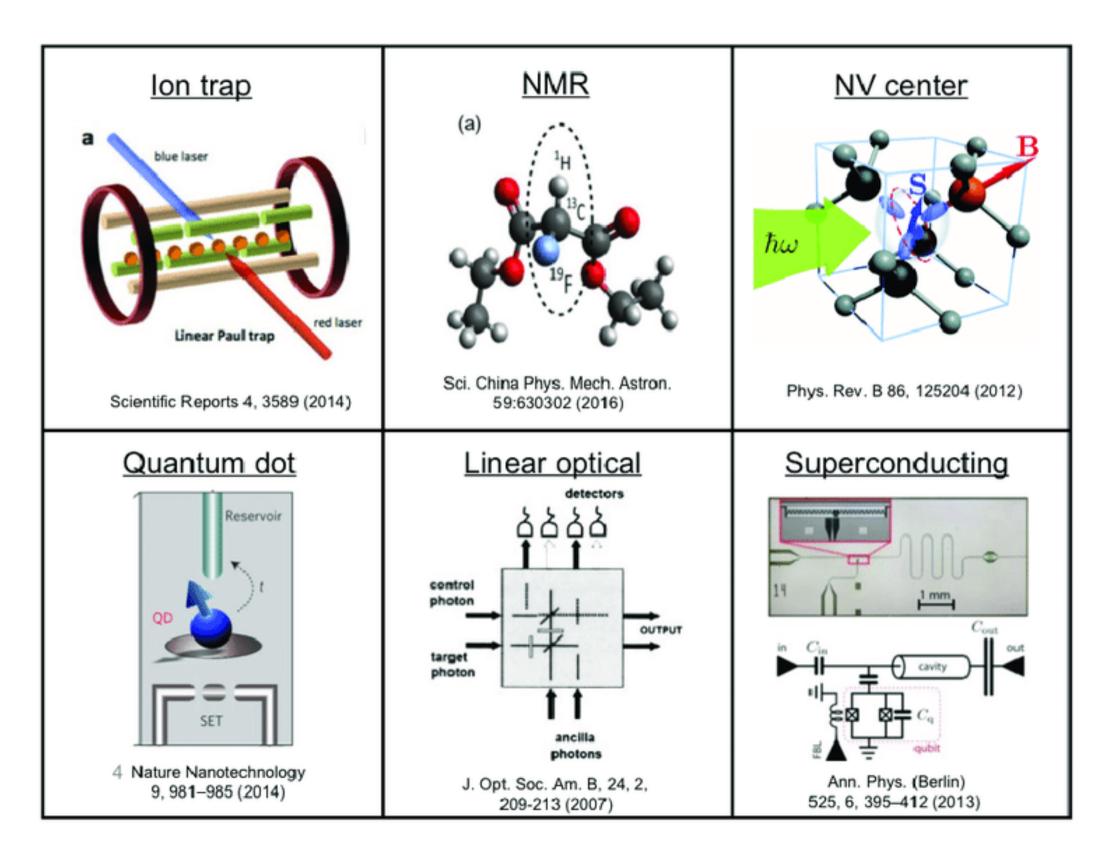
### •Quantum Gate Model

Quantum Annealing

Conclusion

## Many different qubit technologies

• Quantum bit (qubit), the main building block



#### many more

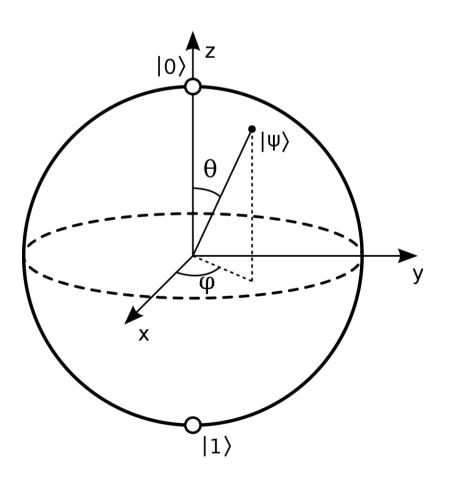
. . .

## What is a qubit (for us)?

• A qubit is a linear combination of basis states

$$|\psi\rangle = \alpha |0\rangle + \beta |1\rangle$$

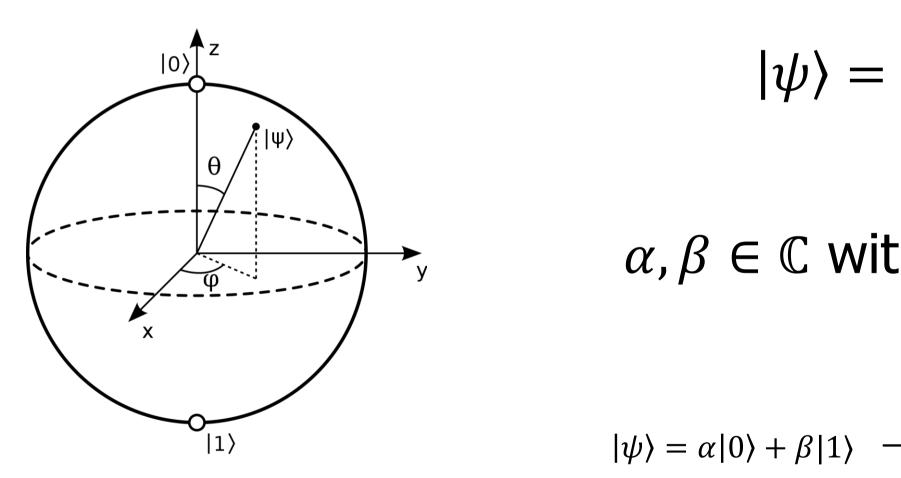
 $\alpha, \beta \in \mathbb{C}$  with  $|\alpha|^2 + |\beta|^2 = 1$ 



22

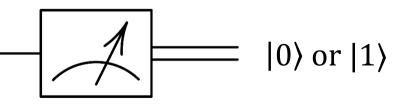
### **Measurement gives classical information**

•  $\alpha$ ,  $\beta$  are called probability amplitudes • When measuring,  $|\alpha|^2$  is probability of finding qubit in state  $|0\rangle$ 



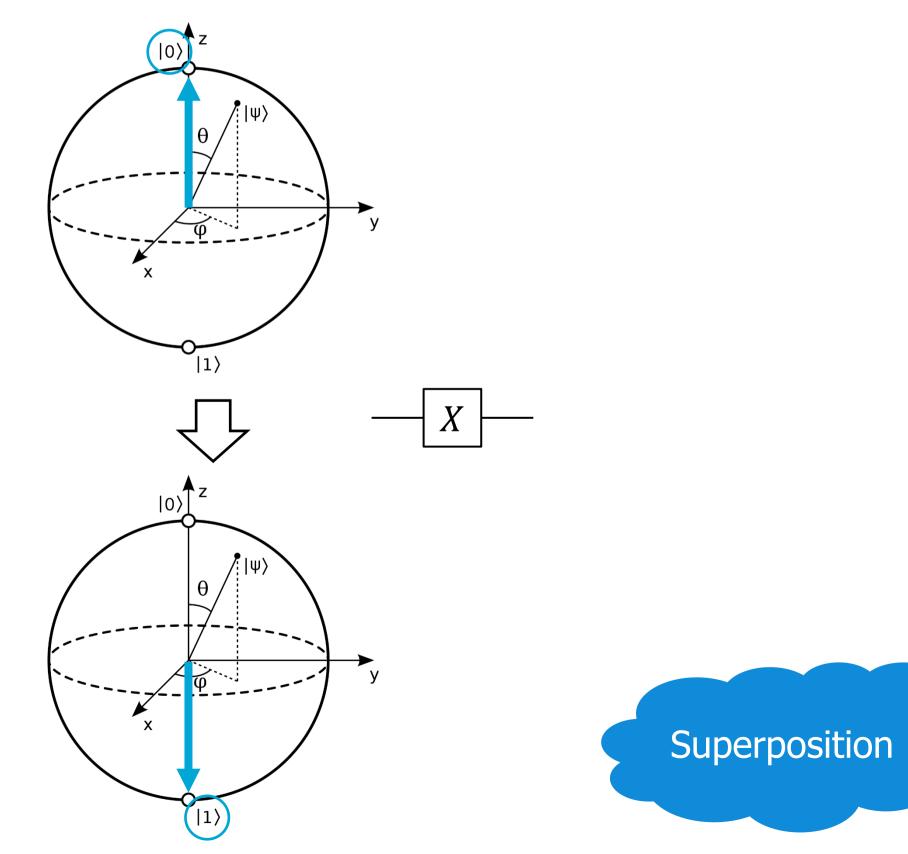
#### $|\psi\rangle = \alpha |0\rangle + \beta |1\rangle$

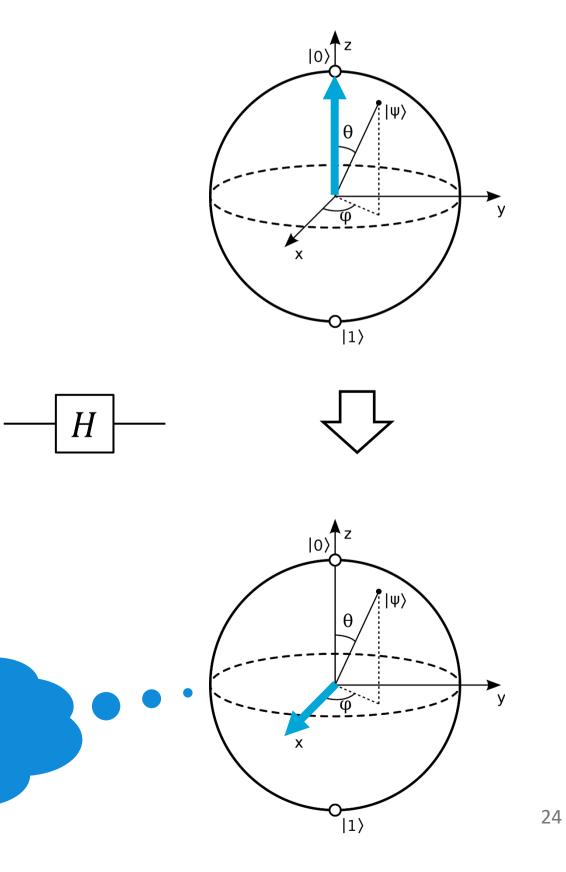
### $\alpha, \beta \in \mathbb{C}$ with $|\alpha|^2 + |\beta|^2 = 1$



### Gates

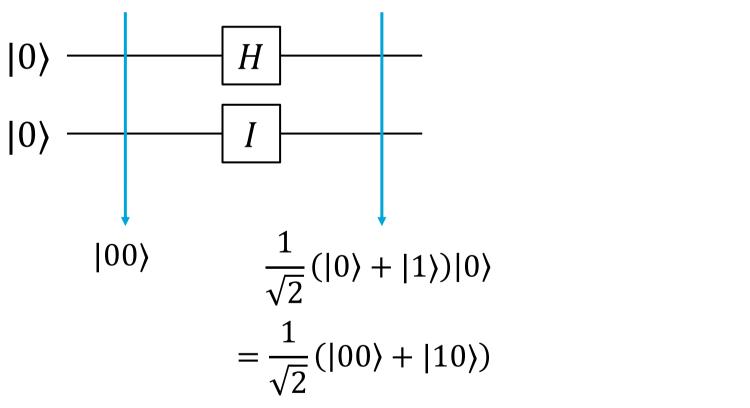
 Quantum gates operate on a quantum state Rotations around an axis, but also controlled operations

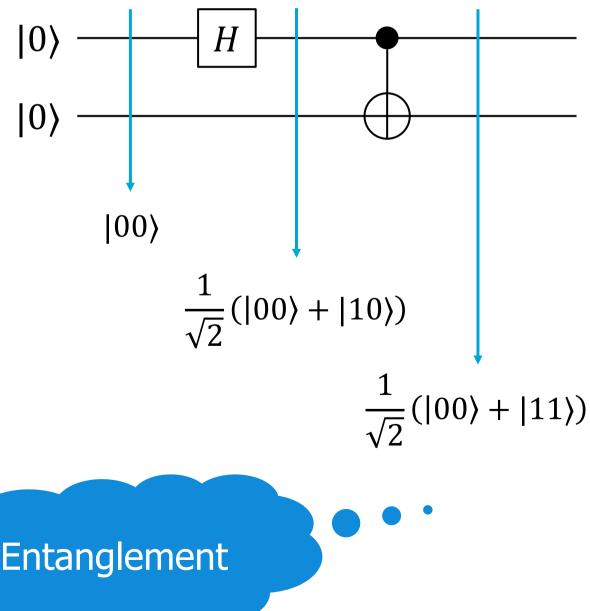


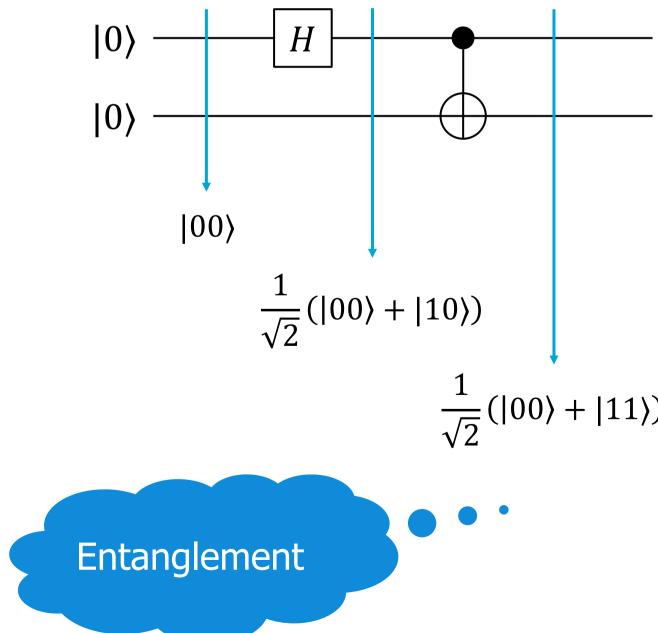


### Gates

 Quantum gates operate on a quantum state Rotations around an axis, but also controlled operations

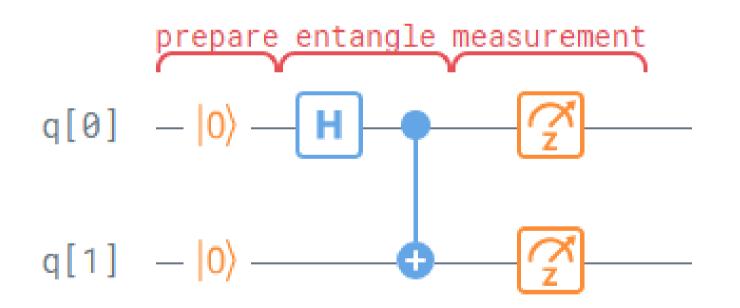






### Circuits

- Back to assembly programming → Quantum Assembly (QASM)
- High-level programming languages available, though

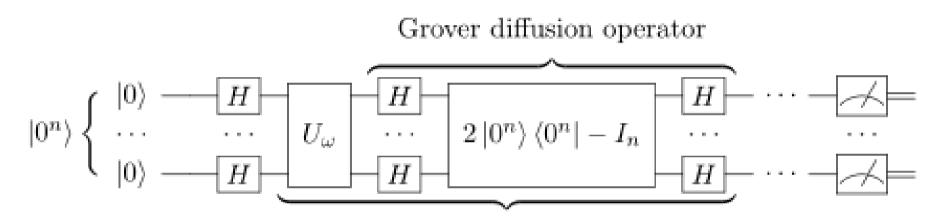


#### tum Assembly (QASM) able, though

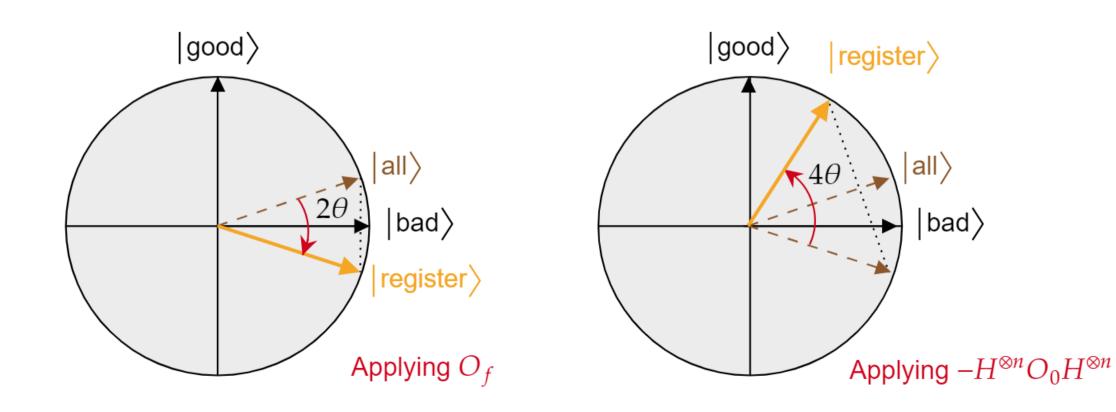
```
version 1.0
 1
2
   qubits 2
 3
 4
    .prepare
 5
        prep_z q[0:1]
 6
 7
    .entangle
 8
        H q[0]
 9
        CNOT q[0], q[1]
10
11
12
    .measurement
        measure_all
13
```

## **General idea of Grover's algorithm**

 Manipulate the amplitudes, such that the probability of our solution state is maximized

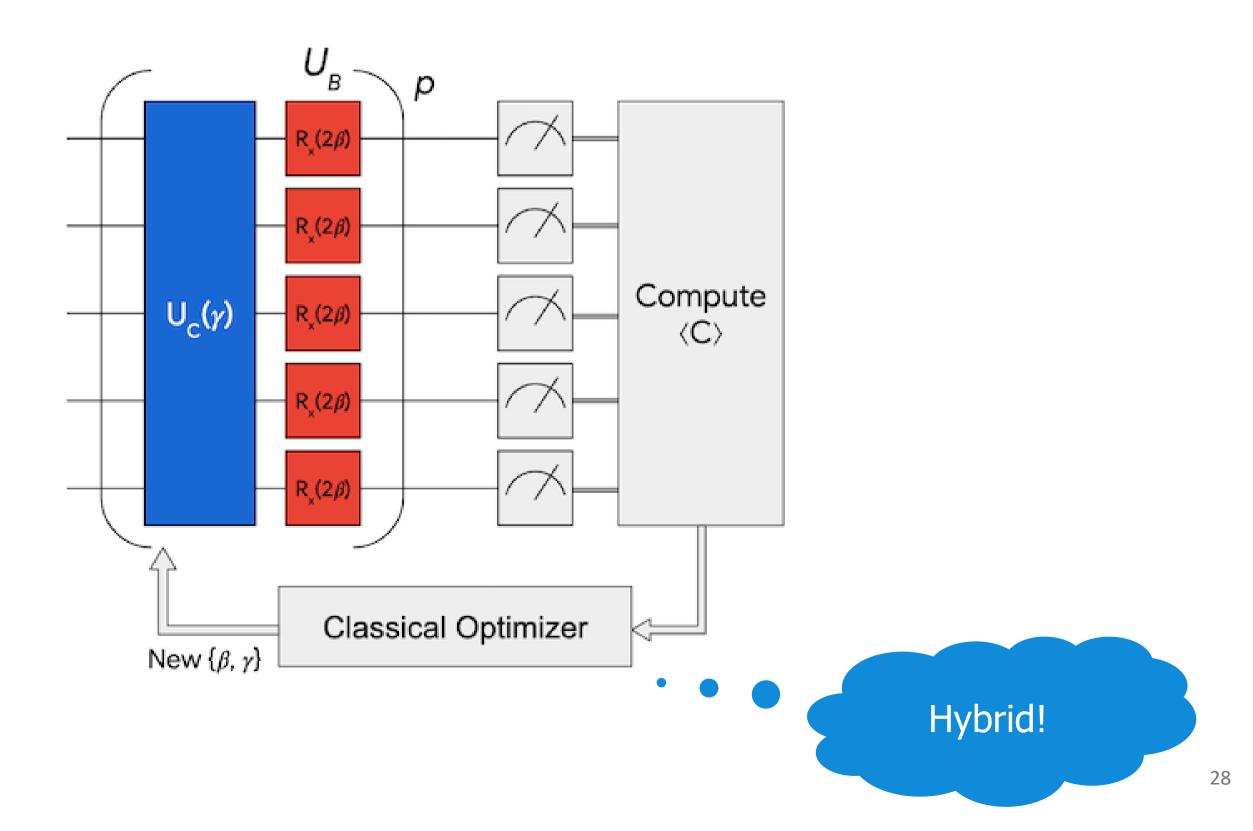


Repeat  $\approx \frac{\pi}{4}\sqrt{N}$  times



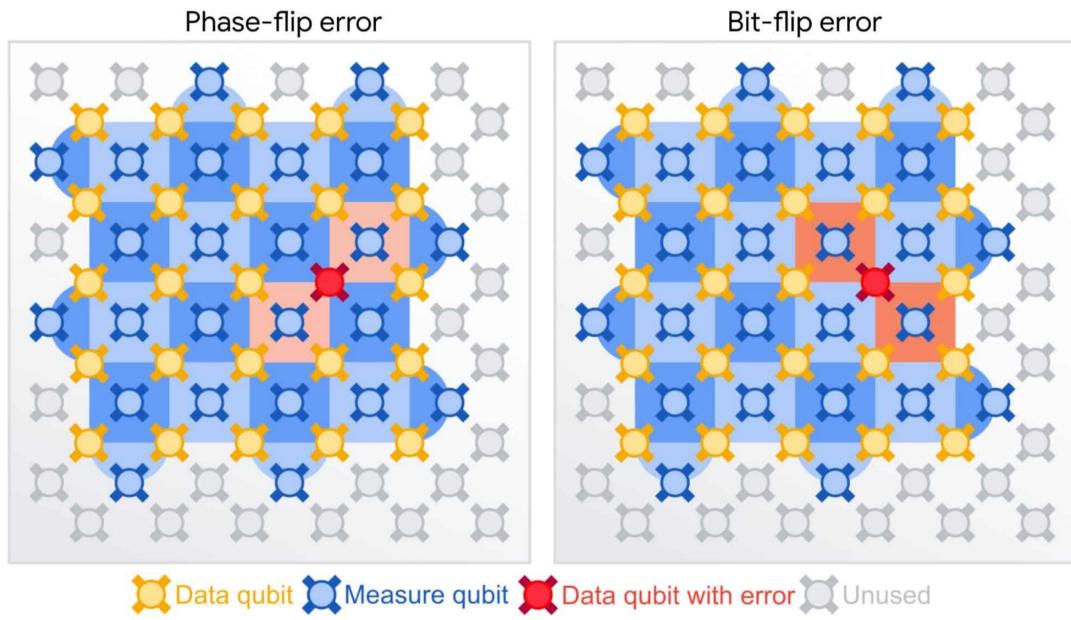
## **General idea of QAOA (and VQE)**

• Minimize a problem-specific quantum operator C by parametrizing  $\beta$  and  $\gamma$ 



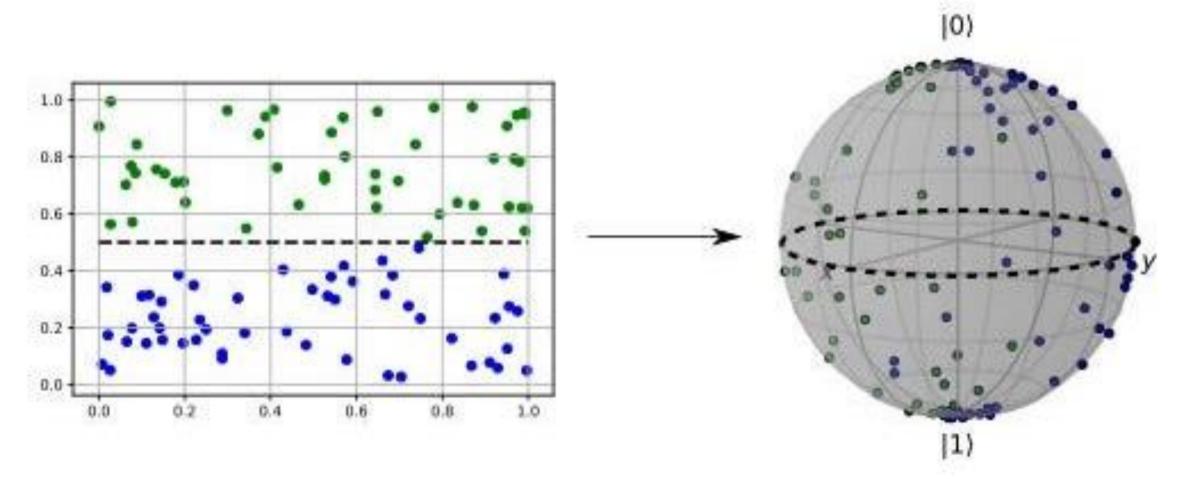
### Main problems

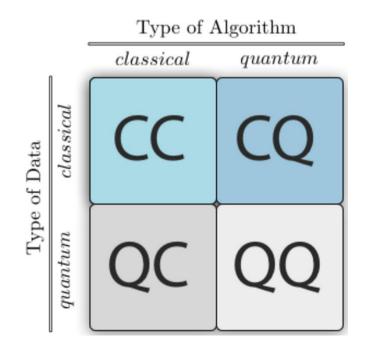
 High noise levels make extensive error correction protocols necessary

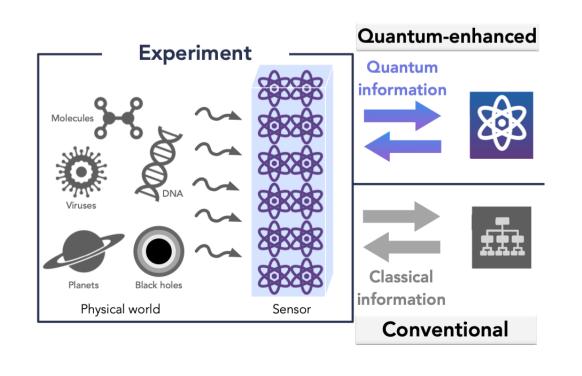


### Main problems

• Data loading, as we basically perform in-memory computing

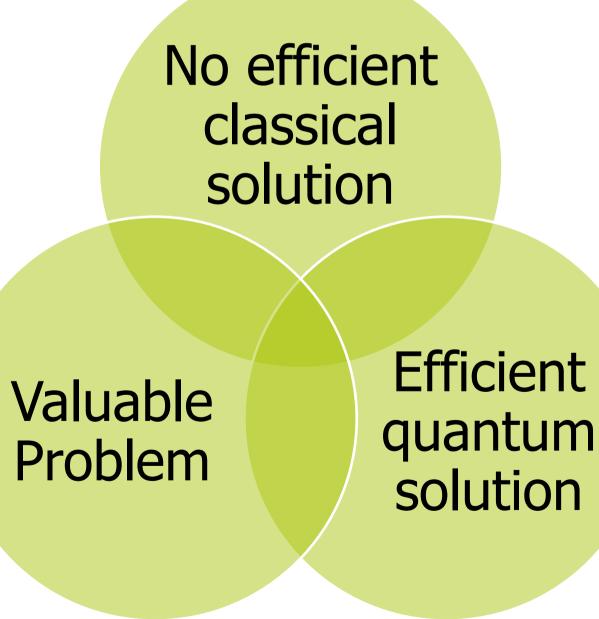






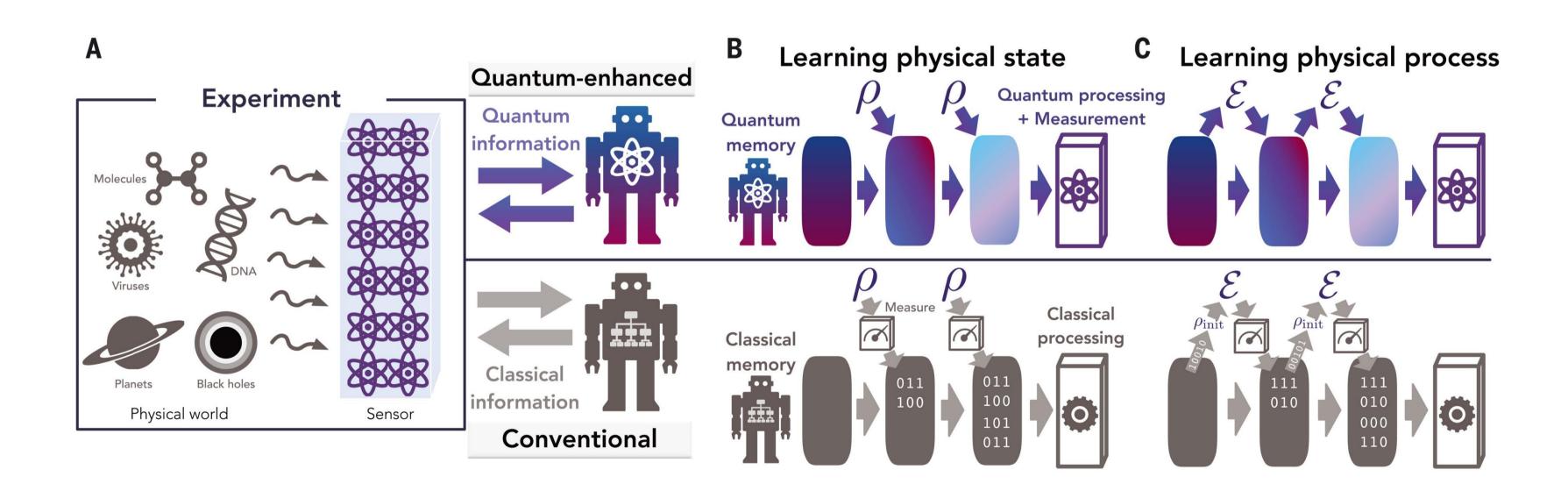
### Main problems

• Find the right solution AND the right problem



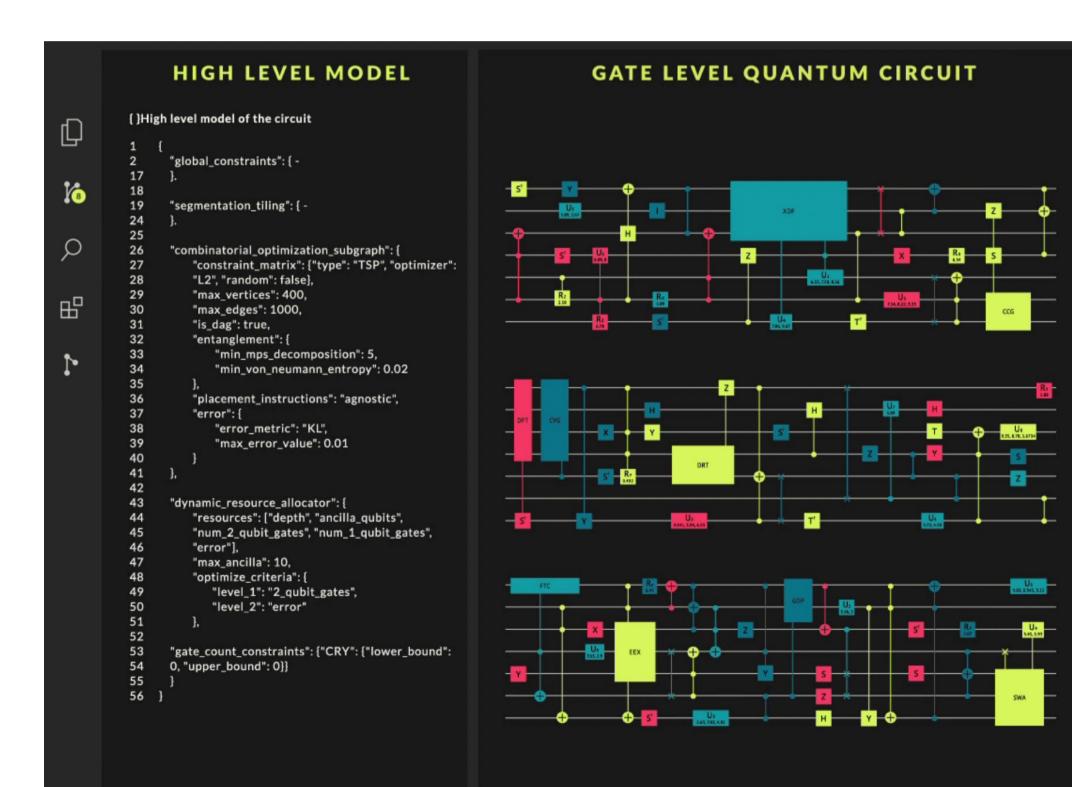
## Some highlights

 Proven superiority, see for example "Quantum advantage in" learning from experiments"



## Some highlights

#### Many successful efforts regarding abstraction, SDKs, APIs, ...



#### GATE LEVEL INSTRUCTIONS

#### //Generated by Classig

| 1  | OPENQASM 2.0;                            |
|----|--|
| 2  | include "qelib1.inc";                    |
| 3  | qreg q[7];                               |
| 4  | h q[4];                                  |
| 5  | ccx q[3],q[0],q[5];                      |
| 6  | ccx q[2],q[1],q[6];                      |
| 7  | u2(6.1795622,3.60187) q[3];              |
| 8  | cswap q[2],q[5],q[4];                    |
| 9  | ccx q[1],q[0],q[6];                      |
| 10 | ch q[1],q[6];                            |
| 11 | rx(6.2680963) q[0];                      |
| 12 | cx q[5],q[4];                            |
| 13 | sdg q[2];                                |
| 14 | sdg q[3];                                |
| 15 | ccx q[1],q[0],q[6];                      |
| 16 | ry(2.0719707) q[2];                      |
| 17 | ccx q[5],q[3],q[4];                      |
| 18 | swap q[3],q[5];                          |
| 19 | ccx q[1],q[6],q[0];                      |
| 20 | cz q[2],q[4];                            |
| 21 | x q[2];                                  |
| 22 | ccx q[6],q[1],q[5];                      |
| 23 | u3(2.8700614,0.43123809,2.8788019) q[0]; |
| 24 | u2(0.72011507,4.9310094) q[4];           |
| 25 | tdg q[3];                                |
| 26 | cz q[0],q[1];                            |
| 27 | cu1(4.4123205) q[5],q[2];                |
| 28 | cz q[4],q[3];                            |
| 29 | t q[6];                                  |
| 30 | rzz(3.6940702) q[2],q[1];                |
| 31 | ccx q[3],q[0],q[4];                      |
| 32 | rx(3.0986421) q[6];                      |
| 33 | sdg q[5];                                |
| 34 | cu1(3.9055758) q[2],q[1];                |
| 35 | cswap q[6],q[5],q[0];                    |
| 36 | tdg q[3];                                |
| 37 | id q[4];                                 |
| 38 | rz(0.57697472) q[0];                     |
| 39 | t q[2];                                  |
| 40 | s q[6];                                  |
| 41 | cu1(5.6409792) q[1],q[4];                |
| 42 | t q[5];                                  |
|    |  |

Overview

Quantum Gate Model

Quantum Annealing

Conclusion

## A niche product in a niche market

 Basically one vendor of quantum annealing • But also: quantum-inspired hardware

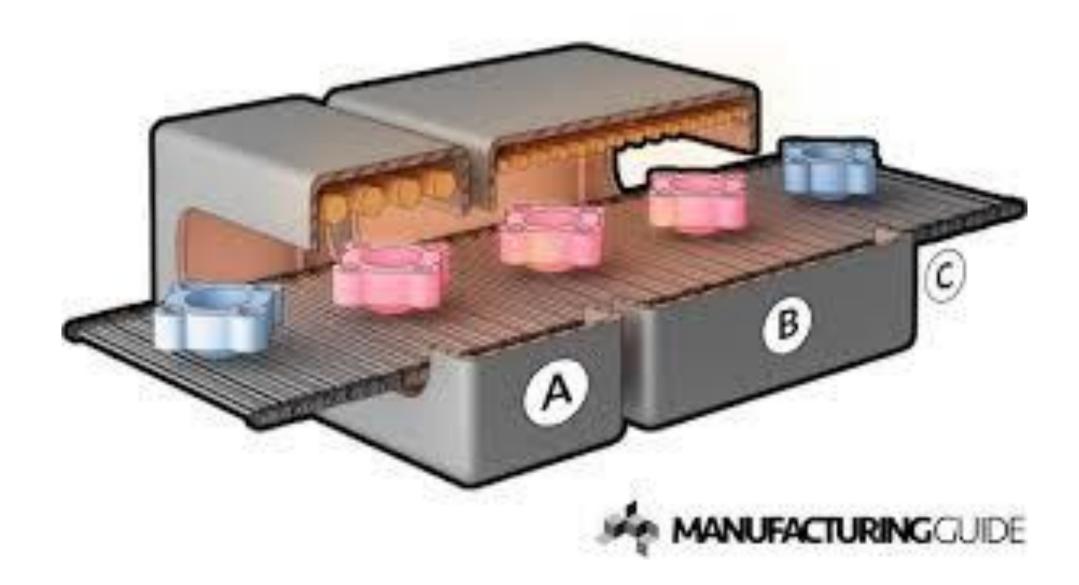




techcrunch.com/2020/09/29/d-wave-launches-its-5000-qubit-advantage-system/ spectrum.ieee.org/fujitsus-cmos-digital-annealer-produces-guantum-computer-speeds

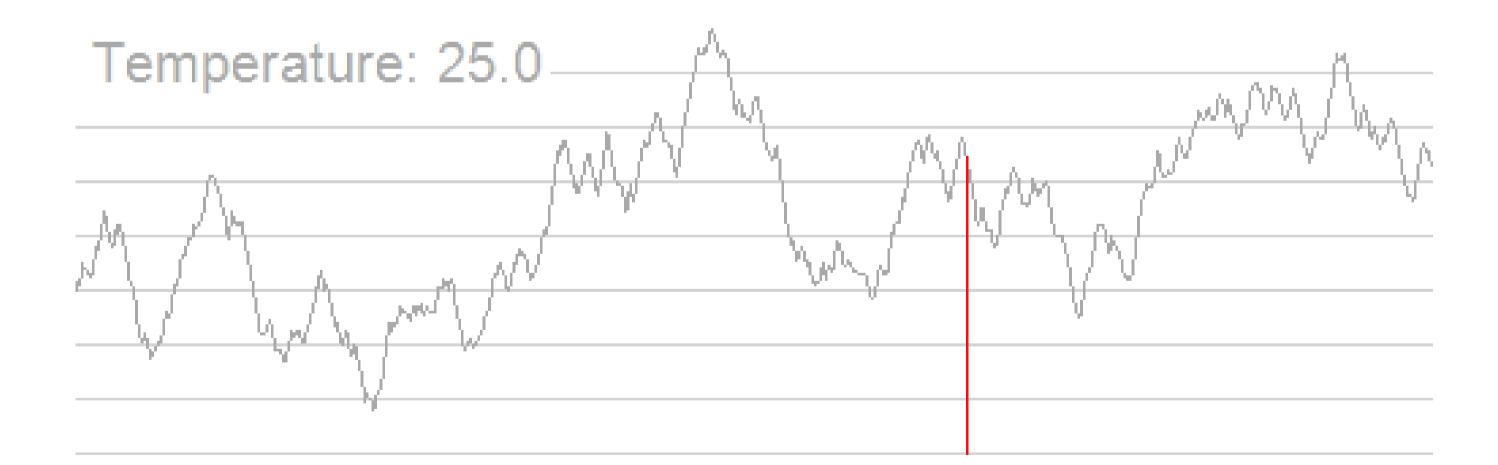
## Annealing in metallurgy and materials science

Heat treatment to alter physical properties of material



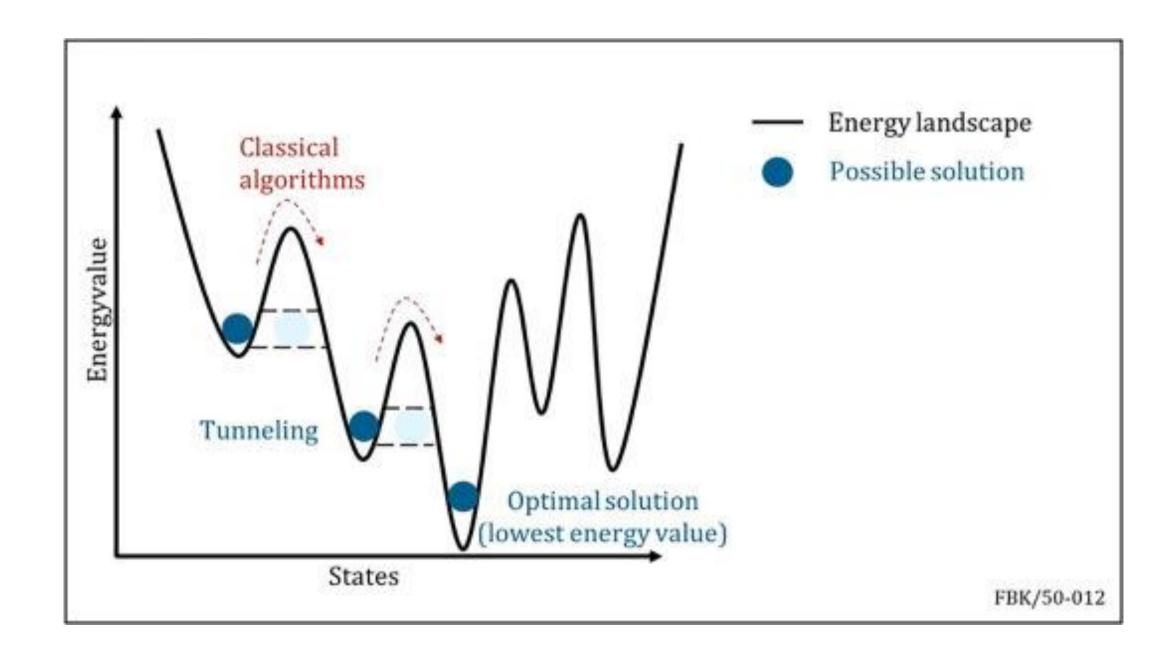
# **Simulated annealing**

- Metaheuristic to find lowest energy state that corresponds to global optimum
- Depending on temperature, worse solution is temporarily accepted



### **Quantum annealing**

Effect of quantum tunneling is utilized



### How to model this solution landscape?

- Hamiltonian function of Ising model from theoretical physics
- Quadratic unconstrained binary optimization (QUBO) from discrete math
- Given a real-valued upper triangular matrix  $Q \in \mathbb{R}^{n \times n}$ • Given a binary vector  $x \in \mathbb{B}^n$
- This defines function  $f_O: \mathbb{B}^n \to \mathbb{R}$  through

$$f_Q(x) = x^T Q x = \sum_{i=1}^n \sum_{j=i}^n Q_i$$

• Goal: find binary vector  $x^*$  that is minimal w.r.t.  $f_0$ , meaning

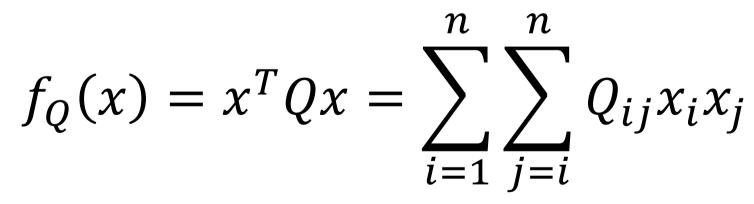
$$\forall x \in \mathbb{B}^n : f_Q(x^*) \le f_Q(x)$$

 $_{i}x_{i}x_{i}$ 

### And now you!

• What binary vector minimizes this QUBO matrix?

| 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 |
|---|---|---|---|---|---|---|---|
|   | 2 | 3 | 4 | 5 | 6 | 7 | 8 |
|   |   | 3 | 4 | 5 | 6 | 7 | 8 |
|   |   |   | 4 | 5 | 6 | 7 | 8 |
|   |   |   |   | 5 | 6 | 7 | 8 |
|   |   |   |   |   | 6 | 7 | 8 |
|   |   |   |   |   |   | 7 | 8 |
|   |   |   |   |   |   |   | 8 |

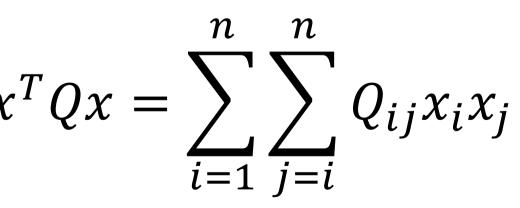


### And now you!

- What binary vector minimizes this QUBO matrix?
- It's f(00000000) = 0
- Any other possibility would add some positive number
- For example: f(00000001) = 8

| 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 |
|---|---|---|---|---|---|---|---|
|   | 2 | 3 | 4 | 5 | 6 | 7 | 8 |
|   |   | 3 | 4 | 5 | 6 | 7 | 8 |
|   |   |   | 4 | 5 | 6 | 7 | 8 |
|   |   |   |   | 5 | 6 | 7 | 8 |
|   |   |   |   |   | 6 | 7 | 8 |
|   |   |   |   |   |   | 7 | 8 |
|   |   |   |   |   |   |   | 8 |

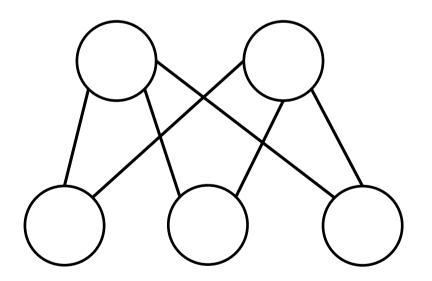
$$f_Q(x) = x$$

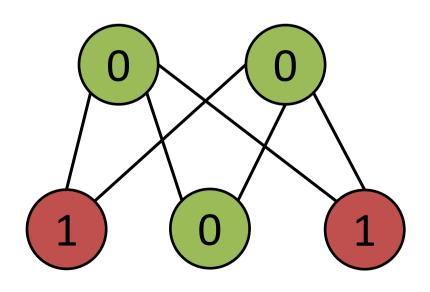


# Maximum cut problem (MAXCUT)

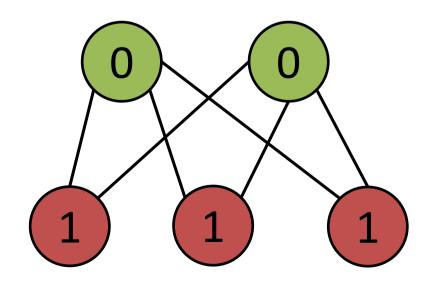
 Goal is to split set of graph's vertices into two disjoint parts ... • ... such that number of edges spanning both parts is maximized

$$\min \sum_{(i,j)\in E} (-x_i - x_j + 2x_i)$$





### $(x_i)$



### Main problems

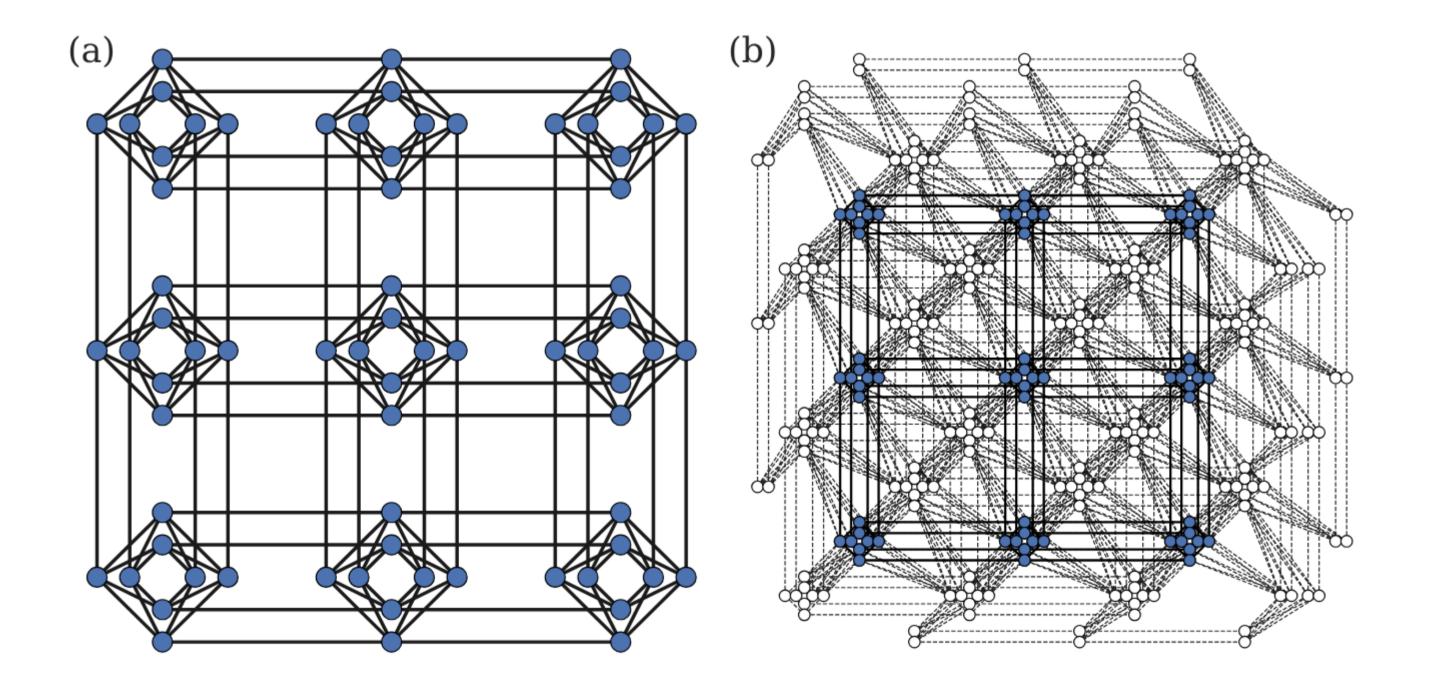
Often inefficient problem transformations

|    | A1       | A2       | A3       | A4       | B1       | B2       | В3       | B4       | C1       | C2       | C3       |
|----|----------|----------|----------|----------|----------|----------|----------|----------|----------|----------|----------|
| A1 | -1<br>-1 | 2        | 2        | 2        | 2        | ab       |          | ab       | 2        | ас       |          |
| A2 |          | -1<br>-1 | 2        | 2        | ab       | 2        | ab       |          | ас       | 2        | ас       |
| А3 |          |          | -1<br>-1 | 2        |          | ab       | 2        | ab       |          | ас       | 2        |
| A4 |          |          |          | -1<br>-1 | ab       |          | ab       | 2        | ас       |          | ас       |
| B1 |          |          |          |          | -1<br>-1 | 2        | 2        | 2        | 2        | bc       |          |
| B2 |          |          |          |          |          | -1<br>-1 | 2        | 2        | bc       | 2        | bc       |
| B3 |          |          |          |          |          |          | -1<br>-1 | 2        |          | bc       | 2        |
| B4 |          |          |          |          |          |          |          | -1<br>-1 | bc       |          | bc       |
| C1 |          |          |          |          |          |          |          |          | -1<br>-1 | 2        | 2        |
| C2 |          |          |          |          |          |          |          |          |          | -1<br>-1 | 2        |
| С3 |          |          |          |          |          |          |          |          |          |          | -1<br>-1 |
| C4 |          |          |          |          |          |          |          |          |          |          |          |

| C4       |
|----------|
| ас       |
|          |
| ас       |
| 2        |
| bc       |
|          |
| bc       |
| 2        |
| 2        |
| 2        |
| 2        |
| -1<br>-1 |
|          |

### Main problems

Minor embedding



### Some highlights

 As with gate model quantum computing, strongest results in the field of(quantum) physics

Article | Published: 19 April 2023

### Quantum critical dynamics in a 5,000-qubit programmable spin glass

Andrew D. King <sup>12</sup>, Jack Raymond, Trevor Lanting, Richard Harris, Alex Zucca, Fabio Altomare, Andrew J. Berkley, Kelly Boothby, Sara Eitemaee, Colin Enderud, Emile Hoskinson, Shuiyuan Huang, Eric Ladizinsky, Allison J. R. MacDonald, Gaelen Marsden, Reza Molavi, Travis Oh, Gabriel Poulin-Lamarre, Mauricio Reis, Chris Rich, Yuki Sato, Nicholas Tsai, Mark Volkmann, Jed D. Whittaker, ... Mohammad H. Amin + Show authors

Nature 617, 61–66 (2023) Cite this article

12k Accesses | 36 Citations | 273 Altmetric | Metrics

### Abstract

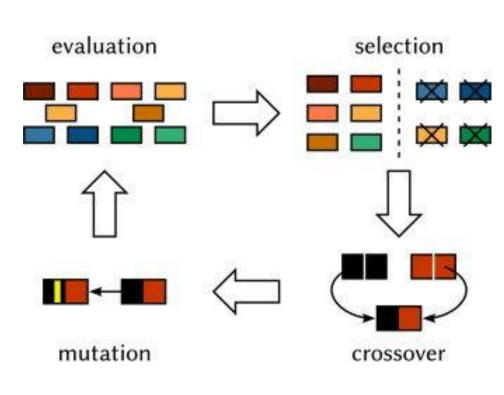
Experiments on disordered alloys1,2,3 suggest that spin glasses can be brought into lowenergy states faster by annealing quantum fluctuations than by conventional thermal annealing. Owing to the importance of spin glasses as a paradigmatic computational testbed, reproducing this phenomenon in a programmable system has remained a central challenge in quantum optimization 4.5.6.7.8.9.10.11.12.13. Here we achieve this goal by realizing quantumcritical spin-glass dynamics on thousands of gubits with a superconducting quantum annealer. We first demonstrate quantitative agreement between quantum annealing and time evolution of the Schrödinger equation in small spin glasses. We then measure dynamics in three-dimensional spin glasses on thousands of qubits, for which classical simulation of many-body quantum dynamics is intractable. We extract critical exponents that clearly distinguish quantum annealing from the slower stochastic dynamics of analogous Monte Carlo algorithms, providing both theoretical and experimental support for large-scale quantum simulation and a scaling advantage in energy optimization.

# Some highlights

- Research on QUBO benefits not only QA community, but also others
  - Gate model (via QAOA)
  - Meta-Heuristics (Genetic/Evolutionary algorithm, Dynamic programming, ...)
  - Operation Research

. . .

Problem transformation, algorithmics, approximation



### Fundamentals of quantum computing

Overview

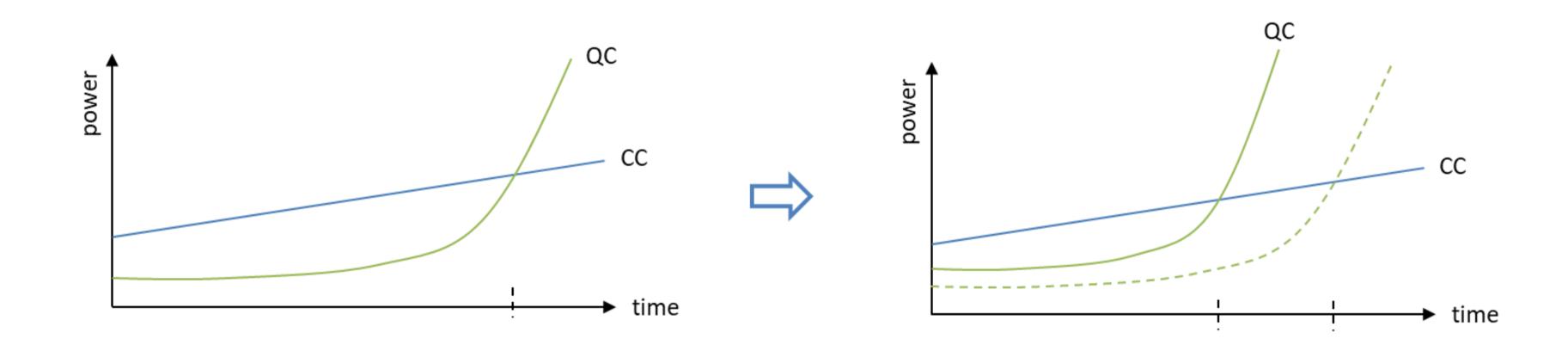
Quantum Gate Model

Quantum Annealing

Conclusion

### Quantum gate model

Become quantum ready!



### Quantum annealing

Find the right problem (transformations)

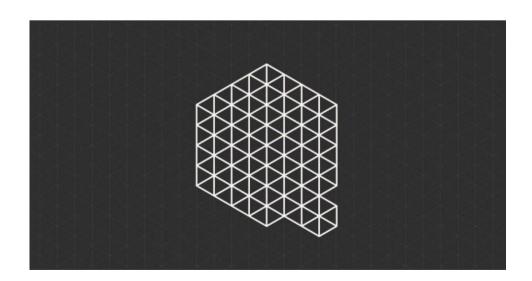
 $\sum^{N} c_{i} X_{i}$ REALO

$$Y_i + \sum_{i=1}^{N} \sum_{j=1}^{i} Q_{ij} X_i X_j$$
$$X_i \in \{0, 1\}$$
$$c_i, Q_{ij} \in \mathbb{R}$$

### Let's get started!

The entrance barrier is lower than one thinks...





... and many more!

# Circ

# 

# Data Management using Quantum Computers -- the Art of Encoding

# LO and Assignment

Learning Objective

Explain how to solve a data management problem on quantum computers

- Assignment
  - Design an approach that solves a data management or data science problem using quantum computers

# **DB Problems Solved Using QPUs**

| Reference   | DB problem                | Subproblem                  | Formulation | Intermediate<br>quantum algorithm | Quantum<br>computer          |
|---|---------------------------|-----------------------------|-------------|-----------------------------------|------------------------------|
| I. Trummer et al.,<br>VLDB'16   | Query optimization        | Multiple query optimization | QUBO        | _                                 | Annealing-based              |
| T. Fankhauser et<br>al., IEEE Access,<br>2023                         |                           |                             |             | QAOA                              | Gate-based                   |
| M. Schonberger et al., SIGMOD23                                       |                           | Join ordering               |             | QAOA                              | Gate-based & annealing-based |
| N. Nayak et al.,<br>BiDEDE '23  |                           |                             |             | QAOA, VQE                         | Gate-based & annealing-based |
| T. Winker et al.,<br>BiDEDE '23                                       |                           |                             | _           | VQC                               | Gate-based                   |
| K. Fritsch et al.,<br>VLDB'23 Demo                                    | Data integration          | Schema matching             | QUBO        | QAOA                              | Gate-based & annealing-based |
| T. Bittner et al.,<br>IDEAS'20, OJCC<br>S. Groppe et al.,<br>IDEAS'21 | Transaction<br>management | Two-phase locking           | QUBO        |                                   | Annealing-based              |

# **Multiple Query Optimization**

- Problem: multiple query optimization (MQO) studies how to choose query plans given a set of queries
  - Goal: minimize the total execution cost
  - Key: shared computation between different queries
  - Valid solution
    - A subset of plans selected for query execution
    - 1:1 mapping between the query and the query plan
  - Optimal solution
    - A valid solution with minimal execution cost among all valid solutions

T. K. Sellis, Multiple-query optimization, TODS 1988



### • Two queries

| Query | Query plan            | Cost |
|-------|-----------------------|------|
| $q_1$ | $p_1$                 | 2    |
|       | $p_2$                 | 4    |
| $q_2$ | $p_3$                 | 3    |
|       | <i>p</i> <sub>4</sub> | 1    |

\* Shared computation saving between  $p_2$  and  $p_3$  is 5

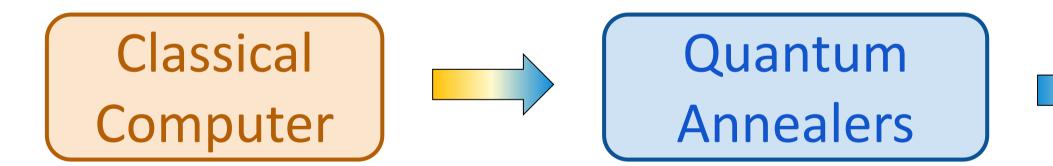
Which query plans to choose for  $q_1$  and  $q_2$ ?

# **Solve MQO Using Quantum Annealing**

Goal: find near-optimal MQO solution

Reduce execution cost by sharing computation among queries

Solution



- Logical mapping 1.
- Physical mapping 2.

Trummer and C. Koch, Multiple query optimization on the D-Wave 2X adiabatic quantum computer," VLDB'16.





### **Encoding: Binary Variables**

**Logical level** 

$$x_p = \begin{cases} 1, \text{ if query plan} \\ 0, \text{ otherwise} \end{cases}$$

**Physical level** 

### *p* is selected

# $b_i = \begin{cases} 1, \text{ if the } i - th \text{ qubit has the state of 1} \\ 0, \text{ if the } i - th \text{ qubit has the state of 0} \end{cases}$

### Example

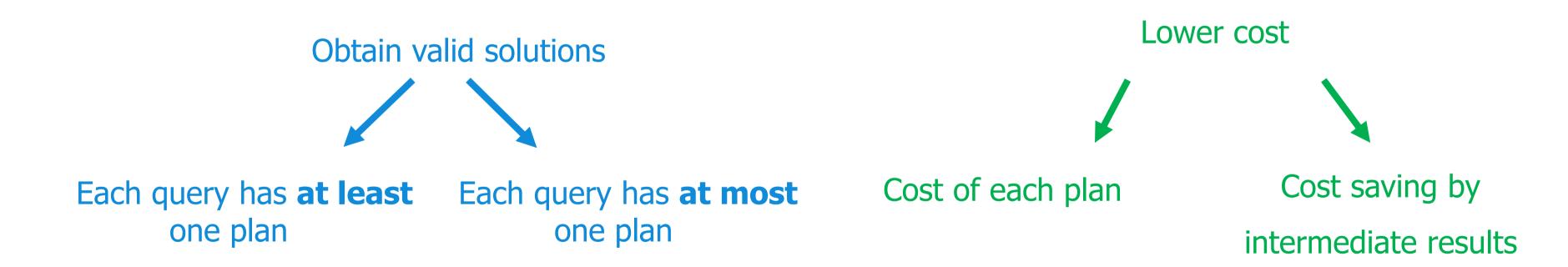
### • Adding binary variable $x_p$

| Query | Query plan            | xp                    | Cost |
|-------|-----------------------|-----------------------|------|
| $q_1$ | $p_1$                 | <i>x</i> <sub>1</sub> | 2    |
|       | <i>p</i> <sub>2</sub> | <i>x</i> <sub>2</sub> | 4    |
| $q_2$ | <i>p</i> <sub>3</sub> | <i>x</i> <sub>3</sub> | 3    |
|       | $p_4$                 | <i>x</i> <sub>4</sub> | 1    |

\* Shared computation saving between  $p_2$  and  $p_3$  is 5

• Design logical mapping:

- Transform the MQO problem into a QUBO problem
- Minimize the logical energy formula







### **1**. Each query has at least one plan

$$E_L = -\sum_{p \in P} x_p$$

### 2. Each query has **at most** one plan

| Query | Query plan | $x_p \in \{0, 1\}$    | Cost |
|-------|------------|-----------------------|------|
| $q_1$ | $p_1$      | <i>x</i> <sub>1</sub> | 2    |
|       | $p_2$      | <i>x</i> <sub>2</sub> | 4    |
| $q_2$ | $p_3$      | <i>x</i> <sub>3</sub> | 3    |
|       | $p_4$      | <i>x</i> <sub>4</sub> | 1    |

\* Shared computation saving between  $p_2$  and  $p_3$  is 5

$$x_1 \rightarrow 1$$
$$x_2 \rightarrow 1$$
$$x_3 \rightarrow 1$$
$$x_4 \rightarrow 1$$

### **NOT** what we want!



1. Each query has at least one plan

$$E_L = -\sum_{p \in P} x_p$$

### 2. Each query has **at most** one plan

$$E_{M} = \sum_{q \in Q} \sum_{\{p_{1}, p_{2} \subseteq Pq\}} x_{p_{1}} x_{p_{2}}$$

| Query | Query plan | $x_p \in \{0, 1\}$    | Cost |
|-------|------------|-----------------------|------|
| $q_1$ | $p_1$      | x <sub>1</sub>        | 2    |
|       | $p_2$      | <i>x</i> <sub>2</sub> | 4    |
| $q_2$ | $p_3$      | <i>x</i> <sub>3</sub> | 3    |
|       | $p_4$      | <i>x</i> <sub>4</sub> | 1    |

\* Shared computation saving between  $p_2$  and  $p_3$  is 5

for  $q_1: x_1 x_2 \rightarrow 0$ for  $q_2: x_3 x_4 \rightarrow 0$ 



### *3. Sum up the cost of each plan*

$$E_C = \sum_{p \in P} c_p x_p$$

If we choose  $p_2$  for  $q_1, p_3$  for  $q_2$ :  $x_1 \rightarrow 0$  $x_2 \rightarrow 1$  $E_C = 0 * 2 + 1 * 4 + 1*3 + 0*1 = 7$  $x_3 \rightarrow 1$  $x_4 \rightarrow 0$ 

| Query | Query plan | $x_p \in \{0, 1\}$    | Cost |
|-------|------------|-----------------------|------|
| $q_1$ | $p_1$      | x <sub>1</sub>        | 2    |
|       | $p_2$      | <i>x</i> <sub>2</sub> | 4    |
| $q_2$ | $p_3$      | <i>x</i> <sub>3</sub> | 3    |
|       | $p_4$      | <i>x</i> <sub>4</sub> | 1    |

\* Shared computation saving between  $p_2$  and  $p_3$  is 5





\* Shared computation saving between  $p_2$  and  $p_3$  is 5

### **4**. Deduct the cost saved by intermediate results

$$E_{S} = -\sum_{\{p_{1}, p_{2} \subseteq P\}} s_{p_{1}, p_{2}} x_{p_{1}} x_{p_{2}}$$

If we choose  $p_2$  for  $q_1, p_3$  for  $q_2$ :

$$x_1 \rightarrow 0$$
  

$$x_2 \rightarrow 1$$
  

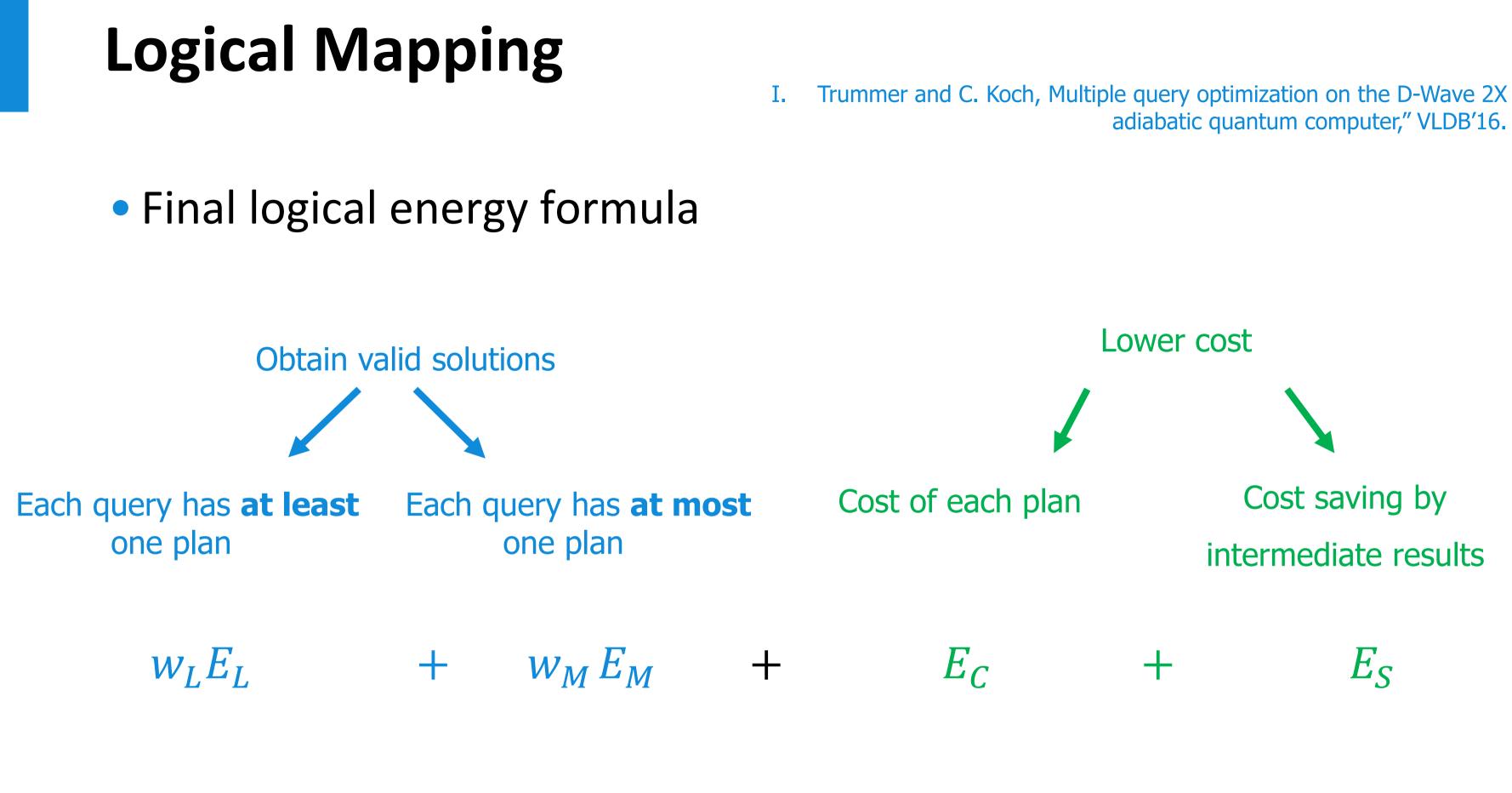
$$x_3 \rightarrow 1$$
  

$$x_4 \rightarrow 0$$
  

$$E_S = -5*1*1 = -$$

| Query | Query plan | $x_p \in \{0, 1\}$        | Cost |
|-------|------------|---------------------------|------|
|       |            | € { <b>0</b> , <b>1</b> } |      |
| $q_1$ | $p_1$      | <i>x</i> <sub>1</sub>     | 2    |
|       | $p_2$      | <i>x</i> <sub>2</sub>     | 4    |
| $q_2$ | $p_3$      | <i>x</i> <sub>3</sub>     | 3    |
|       | $p_4$      | <i>x</i> <sub>4</sub>     | 1    |

### 5



adiabatic quantum computer," VLDB'16.

# **Physical Plan**

- Physical mapping
  - Logical energy formula → physical energy formula of qubit states
  - Key challenge: mapping variables to qubit

- Hardware-specific constraint:
  - Sparse qubit connectivity
  - All qubits representing the same variable form a chain
  - Logical variables in a quadratic term need to be represented by connected groups of qubits
  - Broken qubits

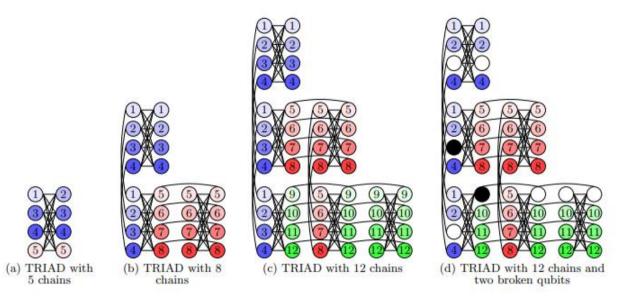


Figure 2: TRIAD pattern in different sizes: we show qubits as circles, annotated by the ID of the logical variable that they represent. The mapping from variables to qubits assures that each variable shares at least one connection (in black) with each of the other variables.

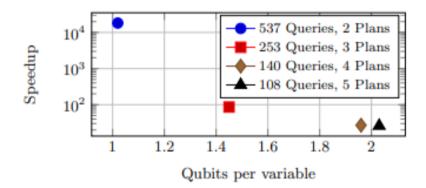
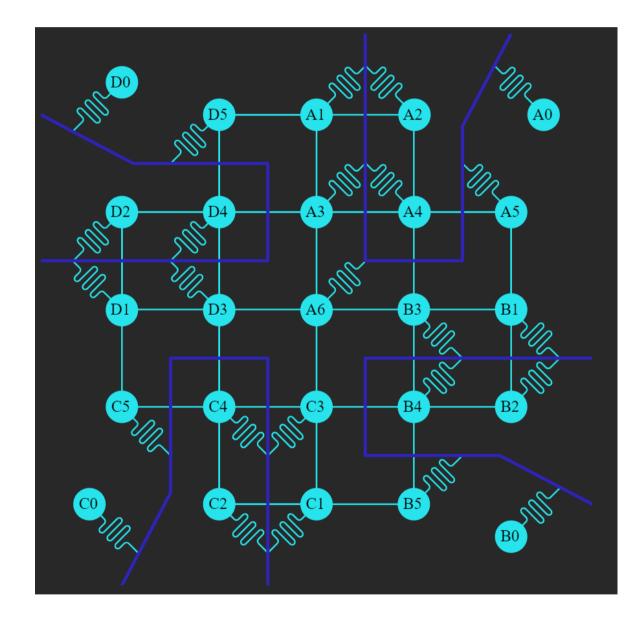


Figure 6: Average quantum speedup for different classes of test cases: having to use more qubits per problem variable decreases the speedup.

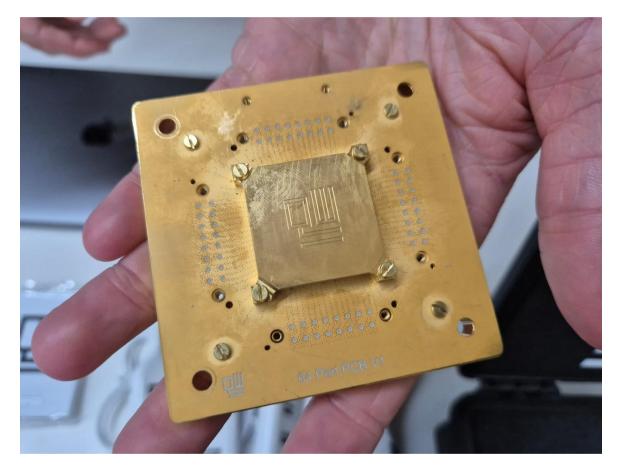
Results: 1000x speedup When fewer qubits per variable

### Hardware

### 25 Qubits: 21 fully connected + 4 isolated qubits







### **Extend MQO to Gate-based QC**

- Gate-based quantum computer, more universal than quantum annealing
- Hybrid classical-quantum algorithms like the Quantum Approximate Optimization Algorithm (QAOA)
- Similar logical mapping, more effective physical mapping
- Limited scalability due to the number of qubits of gate-based QCs

|                      | Trummer and Koch     | This wo    |
|----------------------|----------------------|------------|
| Max. $\#$ of queries | 537 queries, 2 plans | 7 queries, |
| Qubits used          | 1074 / 1097 (98%)    | 14 / 14 (  |
| Max. $\#$ of plans   | 108 queries, 5 plans | 2 queries, |
| Qubits used          | 540 / 1097 (49%)     | 14 / 14 (2 |
|                      |                      |            |

T. Fankhauser, M. E. Soler, R. M. Fuchslin, and K. Stockinger, "Multiple query optimization using a gate-based quantum computer," IEEE Access, 2023.

### ork s, 2 plans (100%), 7 plans (100%)

# Join Ordering

- Problem: join ordering (JO) studies on how to identify the optimal ordering of join operations between relations for an efficient query plan
- Solutions
  - Exhaustive search, e.g., dynamic programming
  - Heuristic methods
  - Special-Purpose Solvers, e.g., MILP
  - ML-based methods
  - SW-HW codesign: methods based on QPU or digital annealer

# Join Ordering on Gate-based QC and QA

- Problem Transformation: from JO to QUBO
- QPU metrics:
  - Overall QPU time
  - Depth of QAOA circuit (gate-based QPUs)
  - Number of qubits

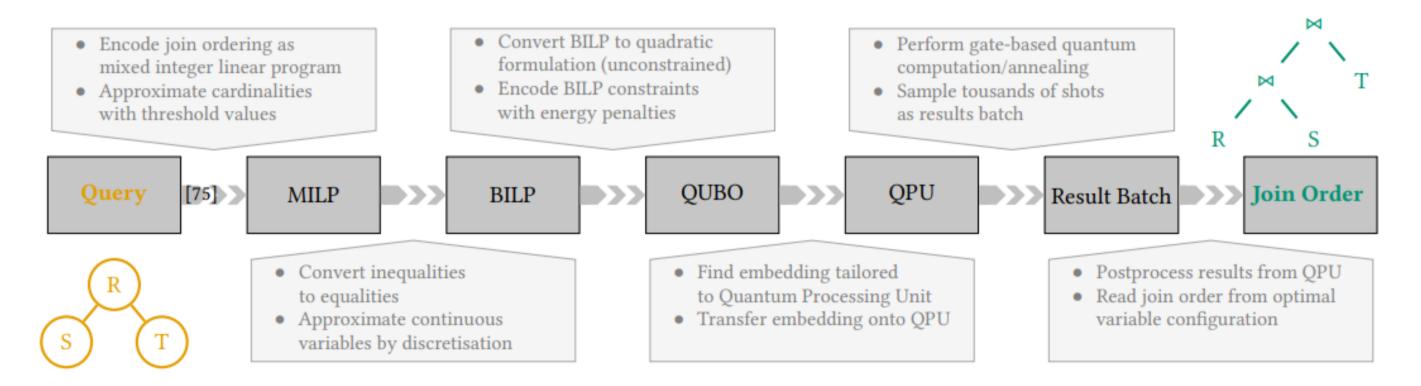


Fig. 1. Overview on all steps required to solve join ordering with quantum computing.

M. Schonberger, et al. "Ready to leap (by co-design)? join order optimisation on quantum hardware. SIGMOD'23.

# Join Ordering using Quantum Annealing

M. Schonberger, et al. "Ready to leap (by co-design)? join order optimisation on quantum hardware. SIGMOD'23.

- Gate-based, QA-based
- Encoding:  $JO \rightarrow MILP \rightarrow QUBO$

Without guaranteed result

M. Schonberger, et al. "Quantum-Inspired Digital Annealing for Join Ordering." VLDB'24.

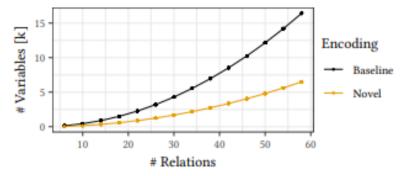


Figure 2: Mandatory variable scaling for cycle queries, comparing the baseline JO-QUBO [41] and our novel encoding.

### Fujitsu 2<sup>nd</sup> Gen digital annealer Fully connected bits

### 8192 bits/variables

### • Encoding: $JO \rightarrow QUBO$

Include parts not supported by MILP, more native to QA/QUBO Efficient with less encoding size With guaranteed result Scalability: 50 relations

### Encoding

### $H = A(HV + H_p) + HC$

Terms:

- A: penalty for violating constraints
- HV: validity term to enforce the solution is valid
- $H_p$ : logarithmic cost of a join
- HC: constraint for intermediate result size

# **More References on Join Ordering Using Quantum Computers**

- 1. M. Schonberger, I. Trummer, and W. Mauerer. "Quantum-Inspired Digital Annealing for Join Ordering." Proceedings of the VLDB Endowment 17.3 (2023): 511-524.
- 2. M. Schonberger, S. Scherzinger, and W. Mauerer, "Ready to leap (by " co-design)? join order optimisation on quantum hardware," Proceedings of the ACM on Management of Data, vol. 1, no. 1, pp. 1–27, 2023.
- M. Schonberger, I. Trummer, and W. Mauerer, "Quantum Optimisation" of General Join Trees," in Joint Workshops 3. at 49th International Conference on Very Large Data Bases (VLDBW'23)—International Workshop on Quantum Data Science and Management (QDSM'23), 2023.
- 4. N. Nayak, J. Rehfeld, T. Winker, B. Warnke, U. C. alikyilmaz, and S. Groppe, "Constructing Optimal Bushy Join Trees by Solving QUBO Problems on Quantum Hardware and Simulators," in Proceedings of the International Workshop on Big Data in Emergent Distributed Environments, ser. BiDEDE '23. 2023.
- T. Winker, U. C. alikyilmaz, L. Gruenwald, and S. Groppe, "Quantum Machine Learning for Join Order Optimization 5. Using Variational Quantum Circuits," in Proceedings of the International Workshop on Big Data in Emergent Distributed Environments, ser. BiDEDE '23, 2023.

### **Schema Matching**

 Schema matching discovers the correspondences among the given source schemas

|   | mortality | pname | age | resting<br>HR |   | mortality | name    | age | oxygen | date<br>diagnosed |
|---|-----------|-------|-----|---------------|---|-----------|---------|-----|--------|-------------------|
| 0 | 0         | Jack  | 20  | 60            | 0 | 1         | Rose    | 45  | 95     | 1/4/21            |
| 1 | 0         | Sam   | 35  | 58            | 1 | 0         | Castiel | 20  | 97     | 3/8/22            |
| 2 | 0         | Ruby  | 22  | 65            | 2 | 1         | Jane    | 37  | 92     | 11/5/21           |
| 3 | 1         | Jane  | 37  | 70            |   | l         |         |     |        |                   |

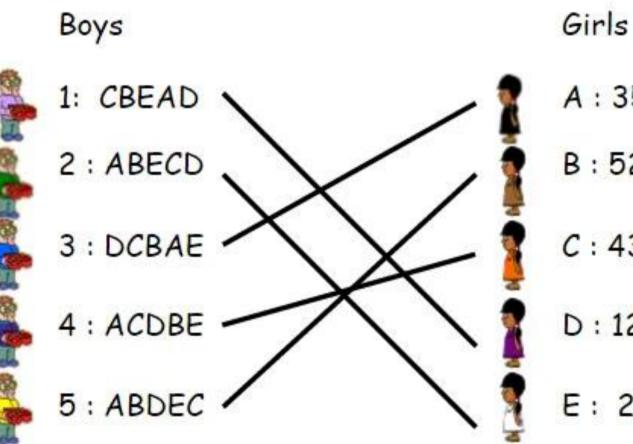
Hospital Amsterdam S<sub>1</sub>

Hospital Delft  $S_2$ 

# **Data Integration -- Schema matching**

Approach

Transform global matching to the stable marriage problem, then into QUBO



An end-to-end hybrid classical and quantum workflow

K. Fritsch and S. Scherzinger, "Solving hard variants of database schema matching on quantum computers," VLDB'23 Demo.

- A: 35214
- B: 52143
- C: 43512
- D:12345
- E: 23415

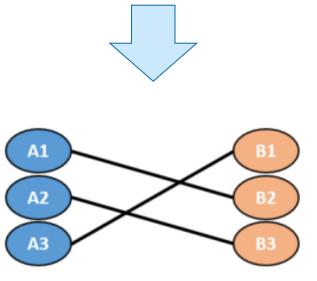
# **Data Integration -- Schema Matching**

 Problem: global matching After obtaining similar attribute pairs with similarity scores, how to assign

one-to-one correspondences, e.g., if there is a tie in ordering

Table 1: Instance of stable matching problem w/ ranked preferences. Brackets denote a tie, i.e., equally preferred partners.

| Set A | Preferences of A <i>i</i> | Set B | Preferences of B <i>i</i> |
|-------|---------------------------|-------|---------------------------|
| A1    | B2, B1                    | B1    | A2, A1, A3                |
| A2    | B3, [B1, B2]              | B2    | A3, A2, A1                |
| A3    | B1, B3, B2                | B3    | A1, A3, A2                |



# **Data Integration -- Schema Matching**

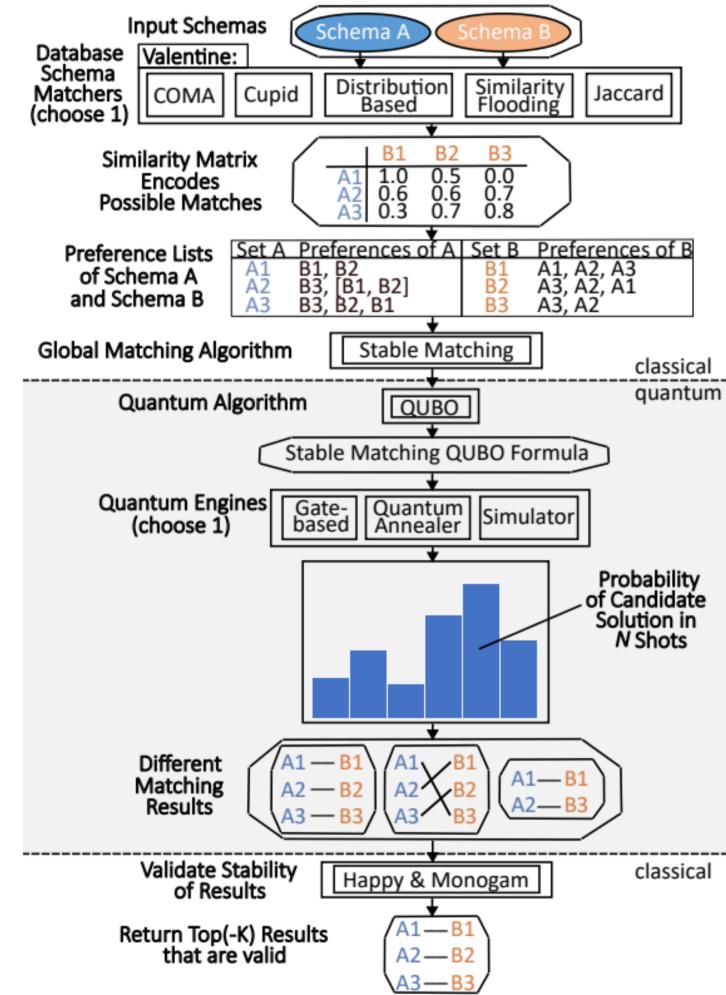
Solution

$$P = -p_1 X_1 + p_2 X_2 + p_3 X_3$$

**Penalties:** 

- $p_1$  : minimize the objective function & to find as many pairs as possible
- $p_2$  : declare for all candidate attribute pairs
- $p_3$ : prevent nonmonogamous matchings (if candidates are matched more than once)

## **Hybrid Classical and Quantum Workflow**





| f | В |
|---|---|
|   |   |
|   |   |

### **Transaction Management**

- Task: distribute n transactions over k machines
- Goal: the distribution is valid and optimal
- Binary variable for QUBO formulation:

$$x_{i, j, s} = \begin{cases} 1, if \ transaction \ t_i \ is \ started \ at \\ 0, otherwise \end{cases}$$

T. Bittner and S. Groppe. 2020. Avoiding Blocking by Scheduling Transactions using Quantum Annealing, IDEAS 2020

### time s on machine m<sub>i</sub>,

### **Transaction Management**

Minimize logical formula

$$P = A + B + C + D$$

### Constraints for a valid schedule:

- A: Each transaction starts exactly once
- B: two or more transactions cannot be executed at the same time on the same machine
- C: transactions that block each other cannot be executed at the
- same time
- Optimal solution
  - D: requires the earliest possible start time for each transaction

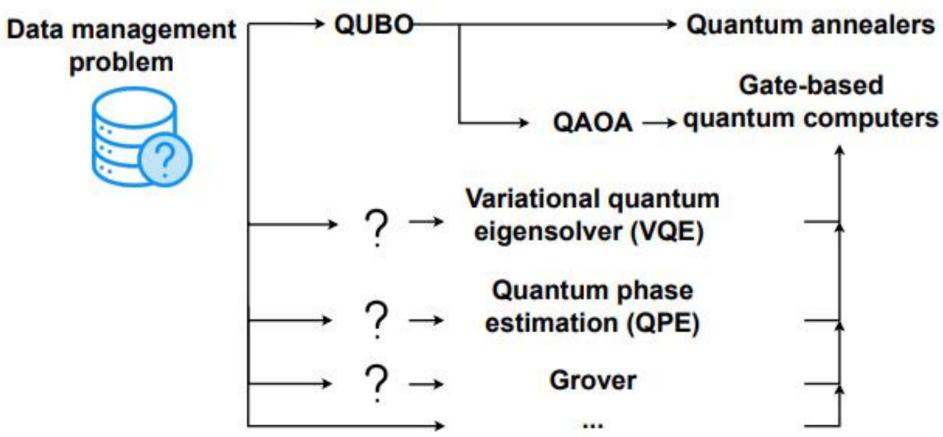
### **More References**

- K. Fritsch and S. Scherzinger, "Solving hard variants of database schema matching on quantum computers," 1. VLDB'23 Demo.
- 2. T. Bittner and S. Groppe, "Avoiding blocking by scheduling transactions using quantum annealing, IDEAS'20.
- 3. Tim Bittner and Sven Groppe, "Hardware accelerating the optimization of transaction schedules via quantum annealing by avoiding blocking, OJCC'20.
- S. Groppe and J. Groppe, "Optimizing transaction schedules on universal quantum computers via code generation 4. for grover's search algorithm, IDEAS'21.
- 5. L. Gruenwald, T. Winker, U. Çalikyilmaz, J. Groppe, and S. Groppe. "Index Tuning with Machine Learning on Quantum Computers for Large-Scale Database Applications." Joint Workshops at 49th International Conference on Very Large Data Bases (VLDBW'23)—International Workshop on Quantum Data Science and Management (QDSM'23). 2023.
- G. Yuan, et al. "Quantum Computing for Databases: A Short Survey and Vision." Joint Workshops at 49th 6. International Conference on Very Large Data Bases (VLDBW'23)—International Workshop on Quantum Data Science and Management (QDSM'23). 2023.

### Roadmap

Solving data management problems on quantum computers

- Problem benefit from quantum advantage, and practically useful
  - Optimization problem
  - Classical approaches have scaling limits
  - Yet it does not require to load a large classical dataset
- Convert a data management solution to quantum algorithms
- Constraints of current quantum hardware

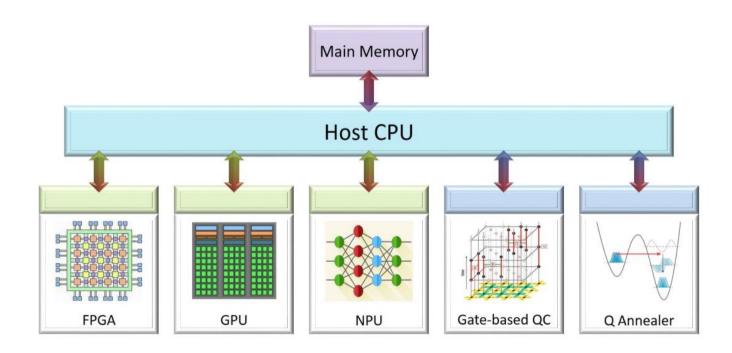




## **Research Opportunities**

- DB problem reformulation
- Hybrid approach on classical and quantum computers
- Optimization given quantum computer constraints





Quantum computer will enhance, not replace, current HPC systems

### m computers <mark>onstraints</mark>



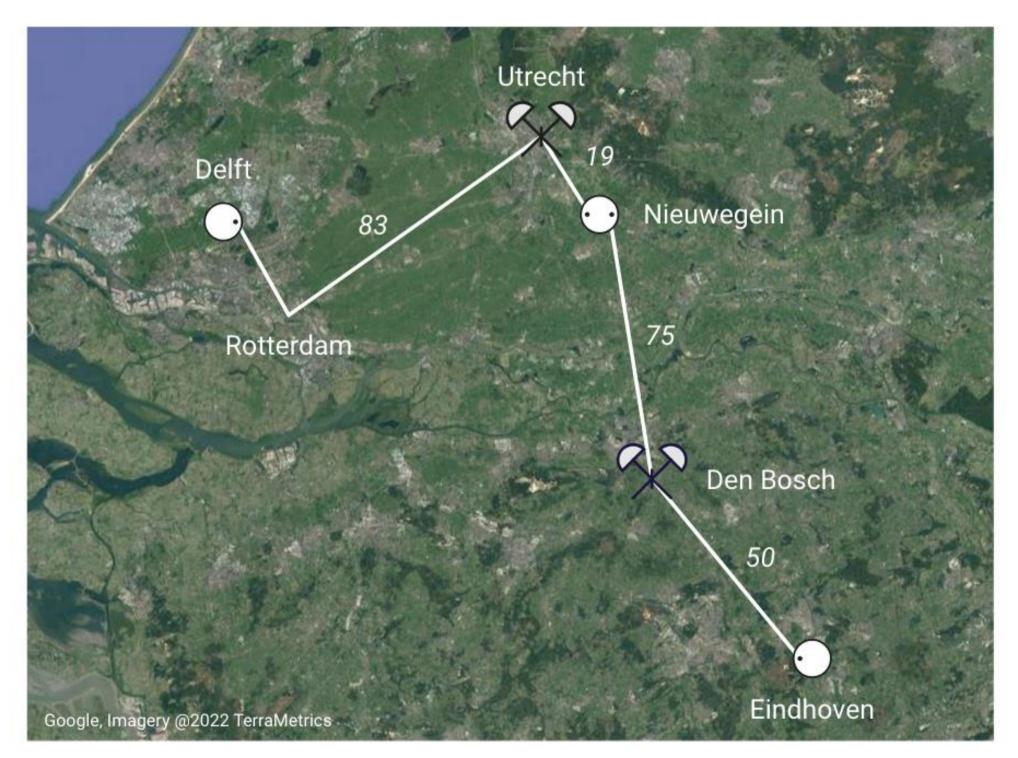
### What is a Quantum Internet?



https://www.youtube.com/watch?v=PCKoT9xcyXI&t=1s

## Hypothetical Quantum Internet Connection

### Total fiber distance between Delft and Eindhoven of 226.5 km

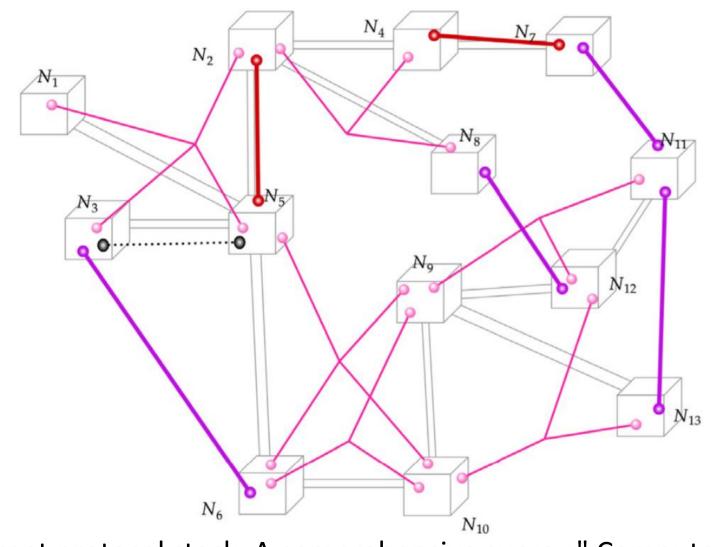


Avis, Guus, et al. "Requirements for a processing-node quantum repeater on a real-world fiber grid." NPJ Quantum

eal-world fiber grid." NPJ Quantum Information 9.1 (2023): 100.

### **Quantum Internet**

- Like classical internet, quantum internet (QI) allows for information exchange between nodes.
  - QI extends CI to allow joint quantum information processing
  - Physical links: channels established to exchange classical messages
  - Virtual links: shared entangled states between the nodes

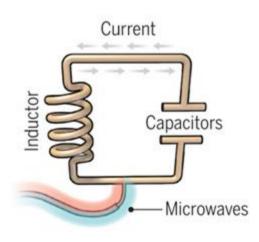


Illiano, Jessica, et al. "Quantum internet protocol stack: A comprehensive survey." Computer Networks 213 (2022): 109092.

## **Different Architectures to Implement Qubits**

### A bit of the action

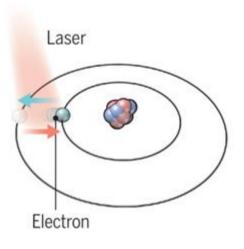
In the race to build a quantum computer, companies are pursuing many types of quantum bits, or gubits, each with its own strengths and weaknesses.





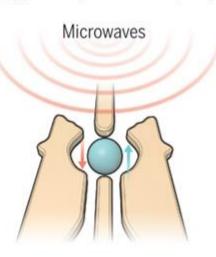
A resistance-free current oscillates back and forth around a circuit loop. An injected microwave signal excites the current into superposition states.

I ongevity (seconds)



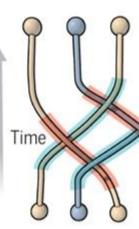
### Trapped ions

Electrically charged atoms, or ions, have quantum energies that depend on the location of electrons. Tuned lasers cool and trap the ions, and put them in superposition states.



### Silicon quantum dots

These "artificial atoms" are made by adding an electron to a small piece of pure silicon. Microwaves control the electron's quantum state.



### **Topological qubits**

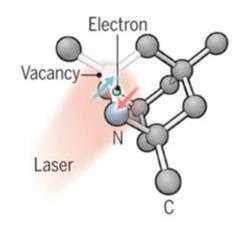
Ouasiparticles can be seen in the behavior of electrons channeled through semiconductor structures. Their braided paths can encode quantum information.

| 0.00005  | >1000  | 0.03  | N/A                     |  |
|--|--|---|-------------------------|--|
| Logic success rate<br>99.4%  | 99.9%  | ~99%  | N/A                     |  |
| Number entangled<br>9  | 14   | 2   | N/A                     |  |
| <b>Company support</b><br>Google, IBM, Quantum Circuits            | ionQ   | Intel   | Microsoft, Bell Labs    |  |
| Pros<br>Fast working. Build on existing<br>semiconductor industry. | Very stable. Highest achieved gate fidelities. | Stable. Build on existing semiconductor industry. | Greatly reduce errors.  |  |
| Collapse easily and must be kept cold.                             | Slow operation. Many lasers are needed.        | Only a few entangled. Must be kept cold.          | Existence not yet confi |  |

Note: Longevity is the record coherence time for a single qubit superposition state, logic success rate is the highest reported gate fidelity for logic operations on two qubits, and number entangled is the maximum number of qubits entangled and capable of performing two-qubit operations.

### Source: https://www.science.org/doi/10.1126/science.354.6316.1090#F2





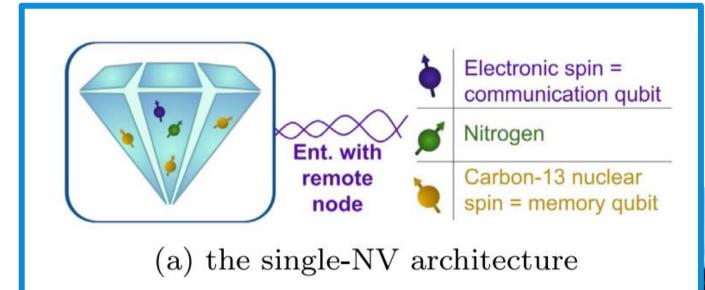
### **Diamond vacancies**

A nitrogen atom and a vacancy add an electron to a diamond lattice. Its quantum spin state, along with those of nearby carbon nuclei, can be controlled with light.

### 10

|         | 99.2%                            |
|---------|----------------------------------|
|         |                                  |
|         | 6                                |
|         | Quantum Diamond Technologies     |
|         | Can operate at room temperature. |
| firmed. | Difficult to entangle.           |

### Quantum Internet Nodes



### Heralded entanglement 1.3km

# Mirrors and filters guide the laser beams to the diamond chip

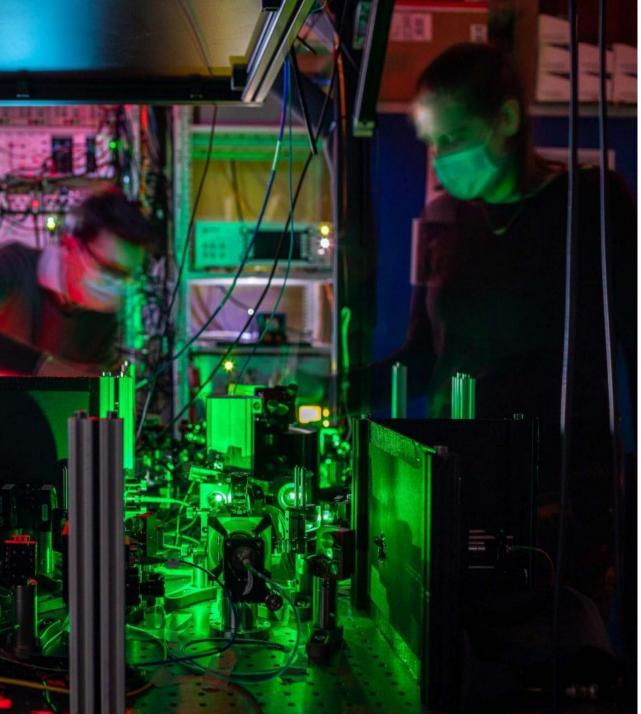
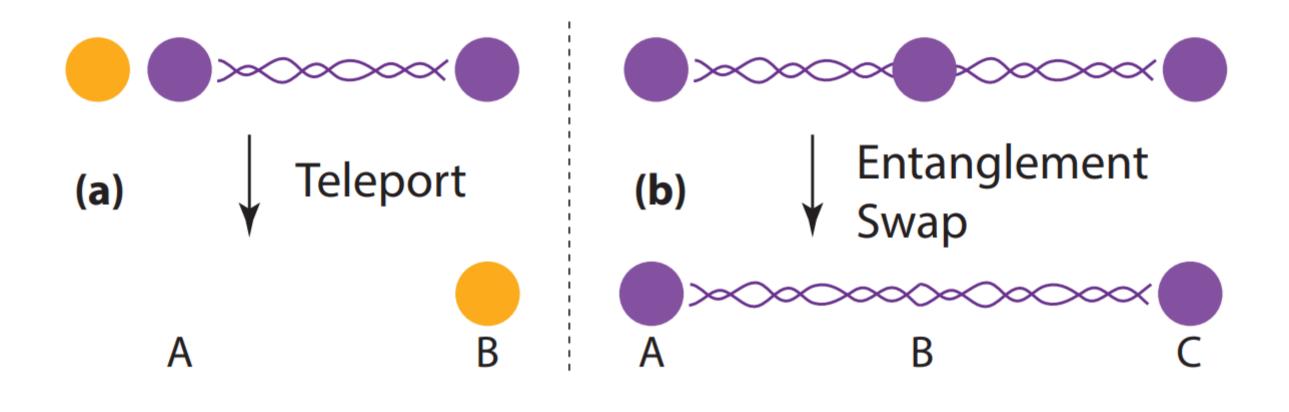


Photo by Marieke de Lorijn for QuTech

96

### **How to Distribute Entanglements?**

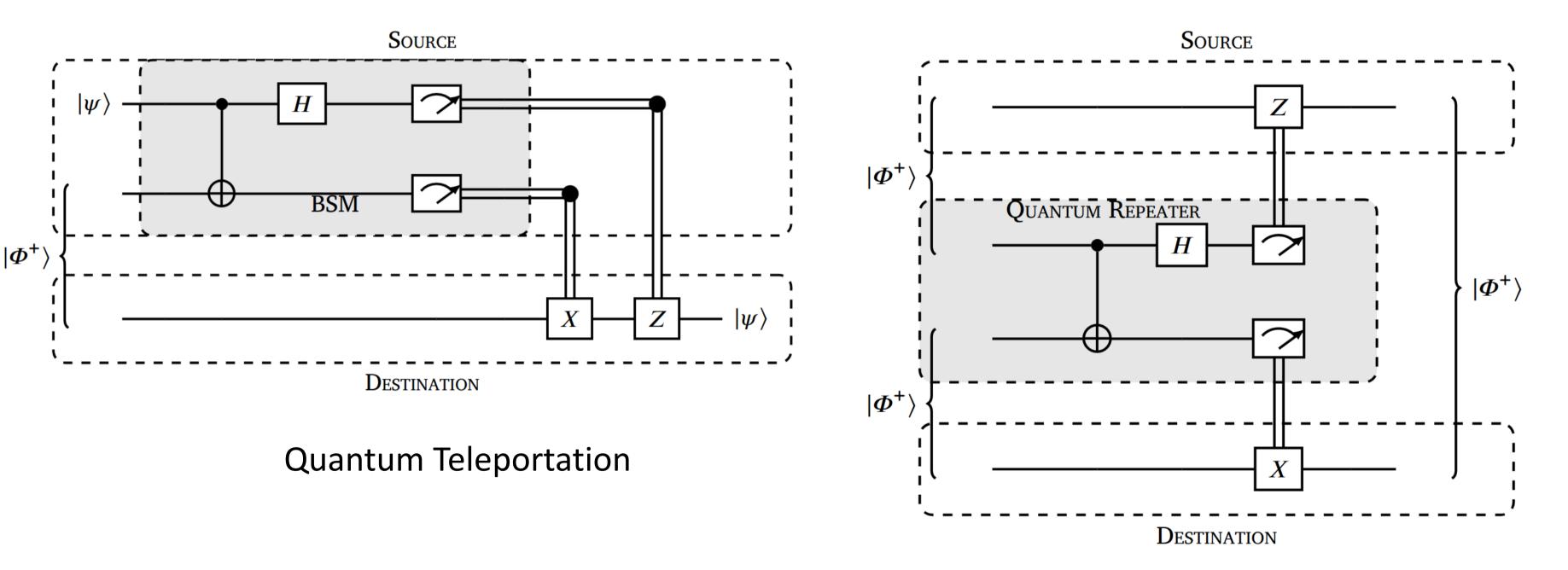


### Quantum Teleportation

Dahlberg, Axel, Matthew Skrzypczyk, Tim Coopmans, Leon Wubben, Filip Rozpędek, Matteo Pompili, Arian Stolk et al. "A link layer protocol for quantum networks." In Proceedings of the ACM special interest group on data communication, pp. 159-173. 2019. 97

### **Entanglement Swapping**

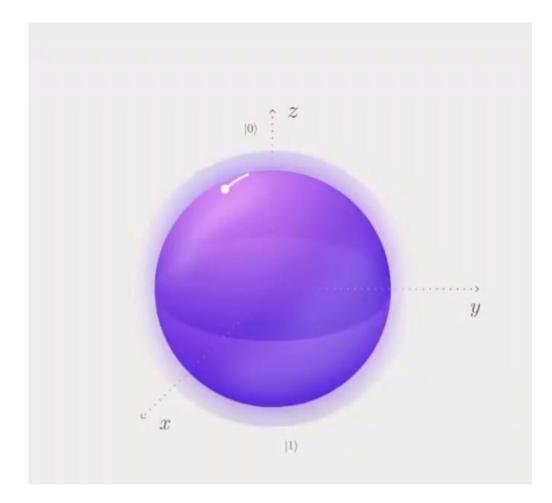
### **How to Distribute Entanglements?**



### **Entanglement Swapping**

## **Quantum Noise**

- Quantum noise makes it hard to extract information from a quantum computer
- Quantum noise results from unwanted coupling with the environment
  - Depolarizing
  - Bit & phase flipping
  - Amplitude & phase damping

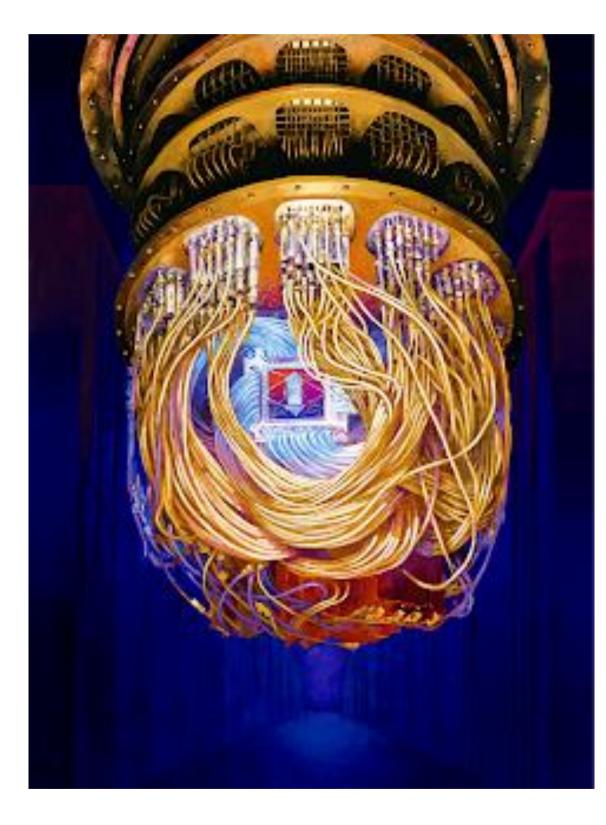




### **Detect Incorrect Quantum Info Processing?**

- We expect to only have noisy intermediate-scale quantum (NISQ) devices in the near future
  - Handle quantum noise by fixing or removing corruption of quantum data?
  - Detect incorrect quantum operations
  - Generate robust entanglements





Google's Sycamore

USTC's Jiuzhang

### What does Quantum Internet Give?

- Full-fledged quantum internet seems necessary for multiparty quantum computation
- Seems a good way to avoid the synchronous problem? Einstein–Podolsky–Rosen (EPR) paradox
- Entanglements enhance data integrity
  - Nonlocal games are useful to detect information leakage and incorrect **operations,** e.g., quantum key distribution (impossible classically)
  - Enhance security of communication impossible by purely classical means

## **Quantum Nonlocality**

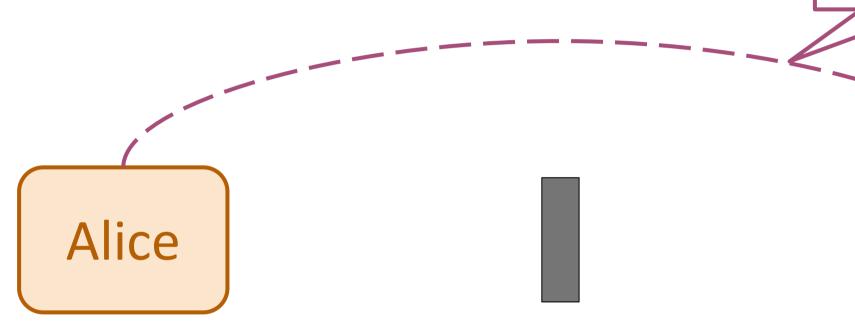
- Correlations that have no interpretations using random coins
- Bell's state: a maximally entangled state that can be used to achieve non-locality

### $|\Psi\rangle = (|00\rangle + |11\rangle)/\sqrt{2}$

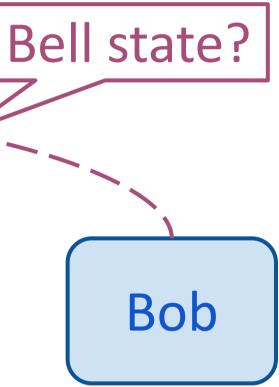
Measuring any half of the state in any basis, the other "collapses" to the same outcome.

 Nonlocal games: Tests used to test nonlocality, specifically, the shared entanglements!

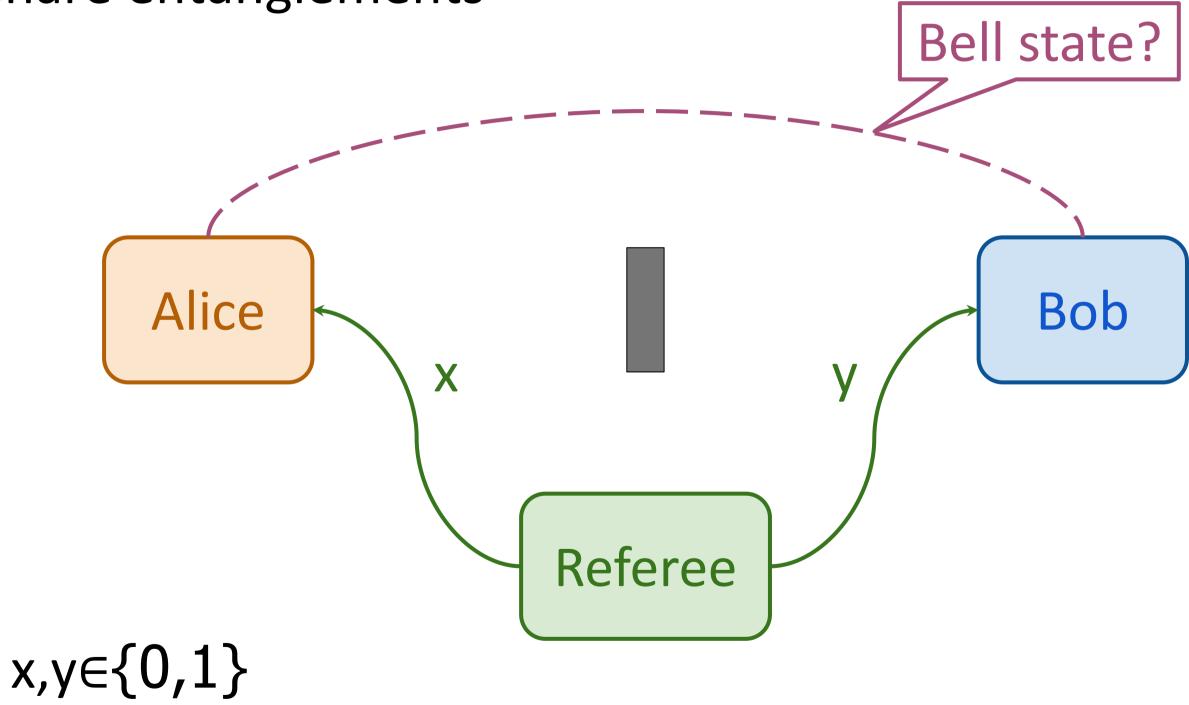
 Two devices that are not allowed to communicate, but can share entanglements



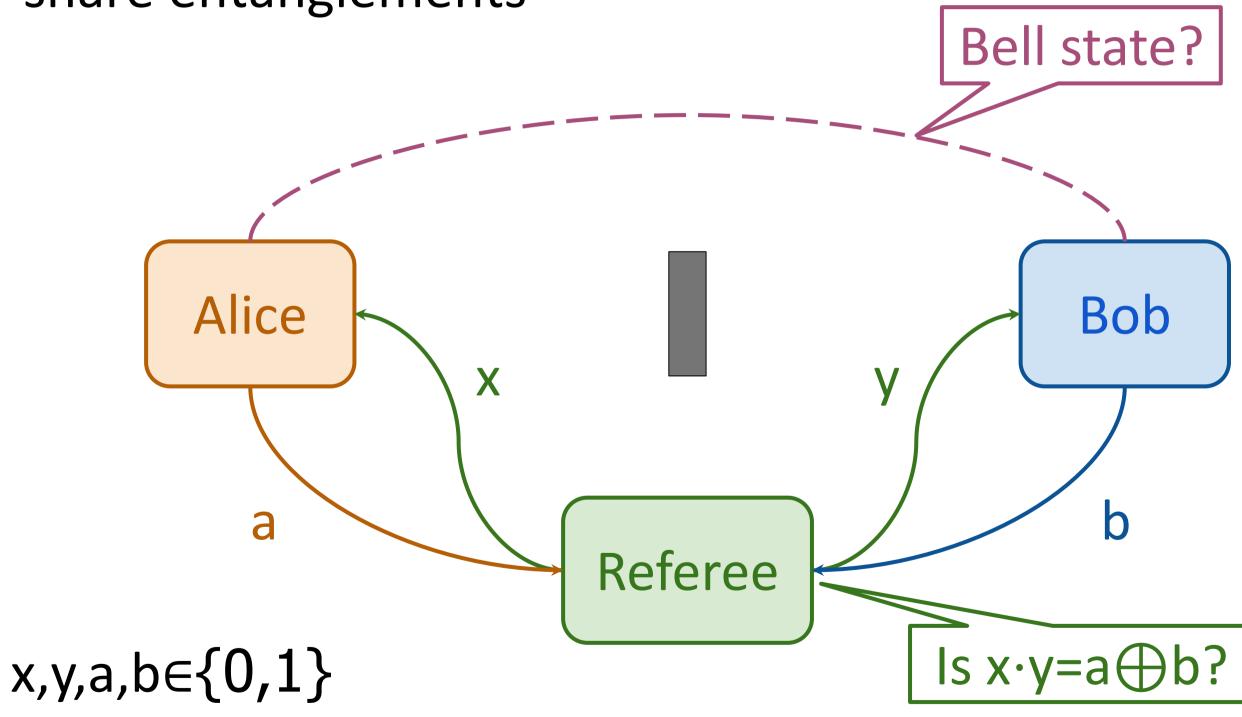




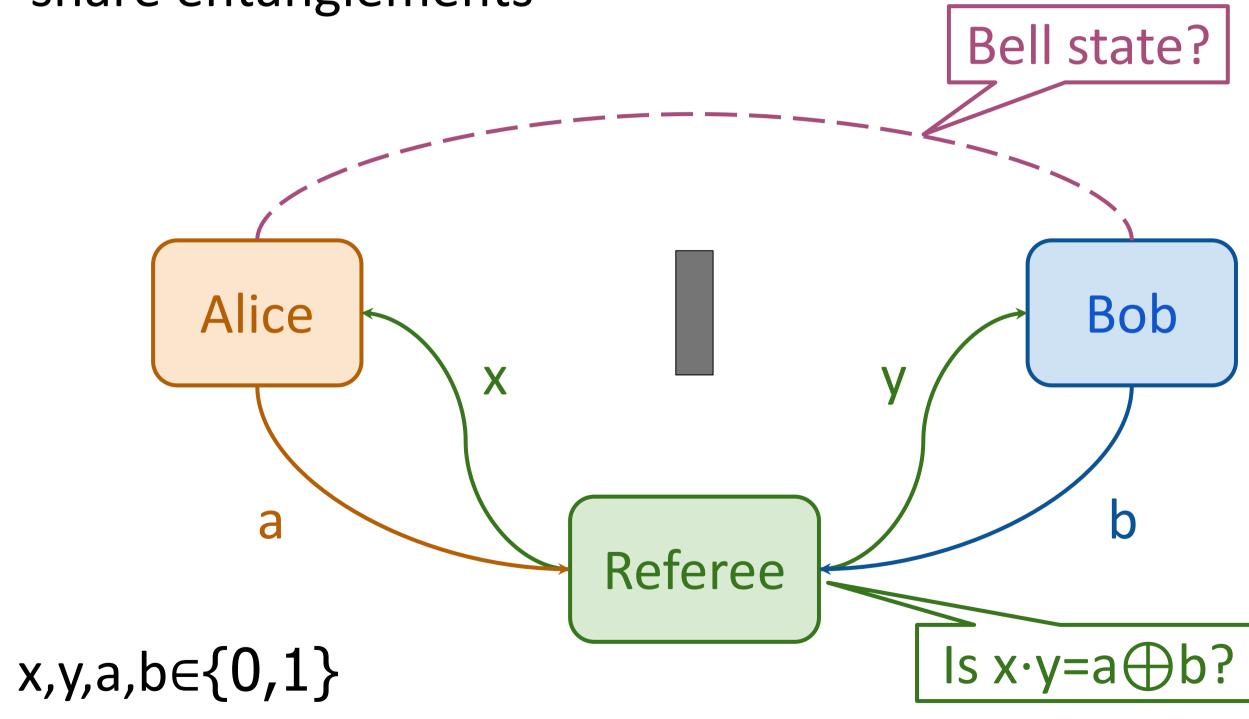
 Two devices that are not allowed to communicate, but can share entanglements



 Two devices that are not allowed to communicate, but can share entanglements



Two devices that are not allowed to communicate, but can share entanglements

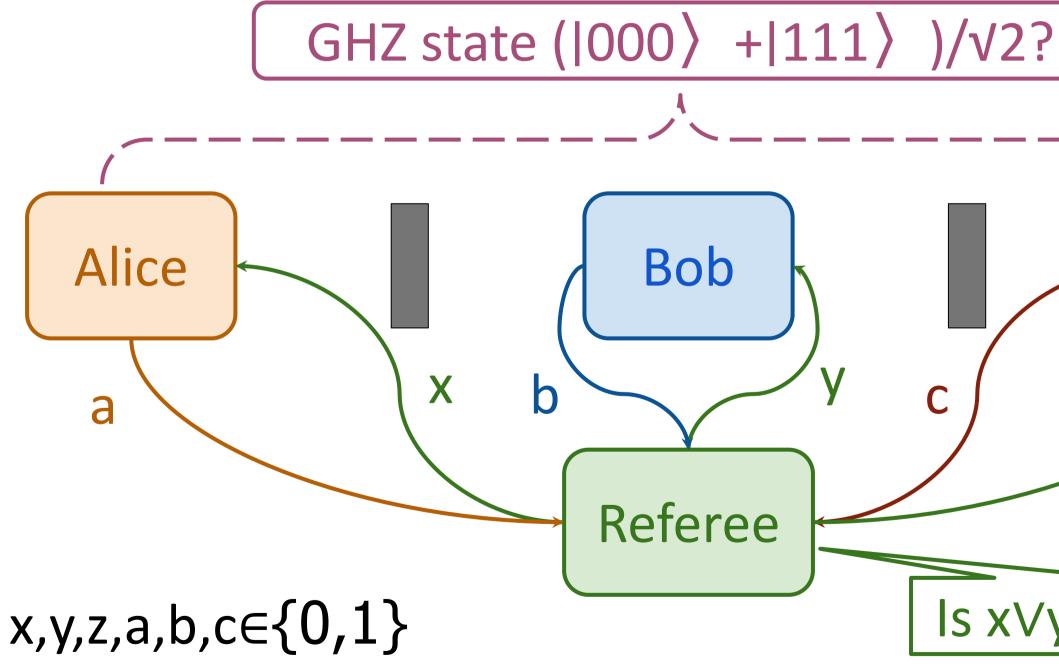


With a Bell state, they can win the test with probability  $\approx 0.85!$ 

107

# The Greenberger-Horne-Zeilinger (GHZ) Game

- Extension to test multipartite entanglements
- A three-player non-local game for testing GHZ state



With a GHZ state, they can win the test with probability 1!

- Charlie Is  $xVyVz=a\oplus b\oplus c$ ?

# **Beyond the Binary Alphabet**

For data management, it is essential to consider the space of data beyond binary strings.

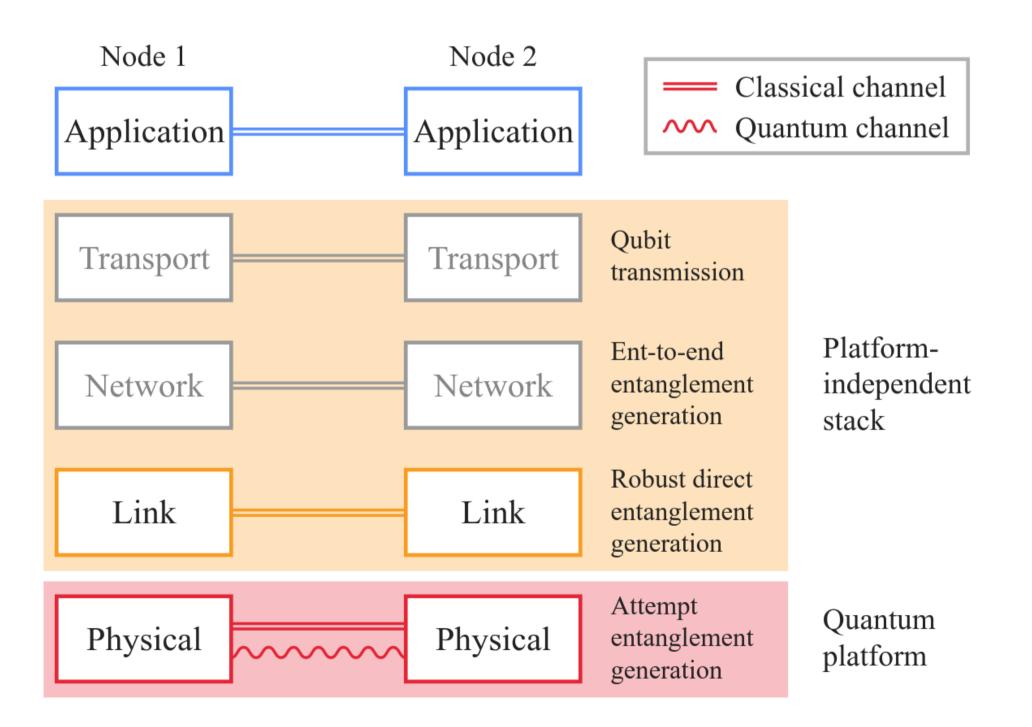
The tests may be generalized to "qudits," which takes values from an alphabet of size d>2 (as opposed to  $\{0,1\}$ ).

| pname   | gender | age | restin<br>HR |
|---------|--------|-----|--------------|
| Alice   | 0      | 96  | 60           |
| Bob     | 1      | 35  | 70           |
| Charlie | 1      | 22  | 65           |

### ng mortality 0 1

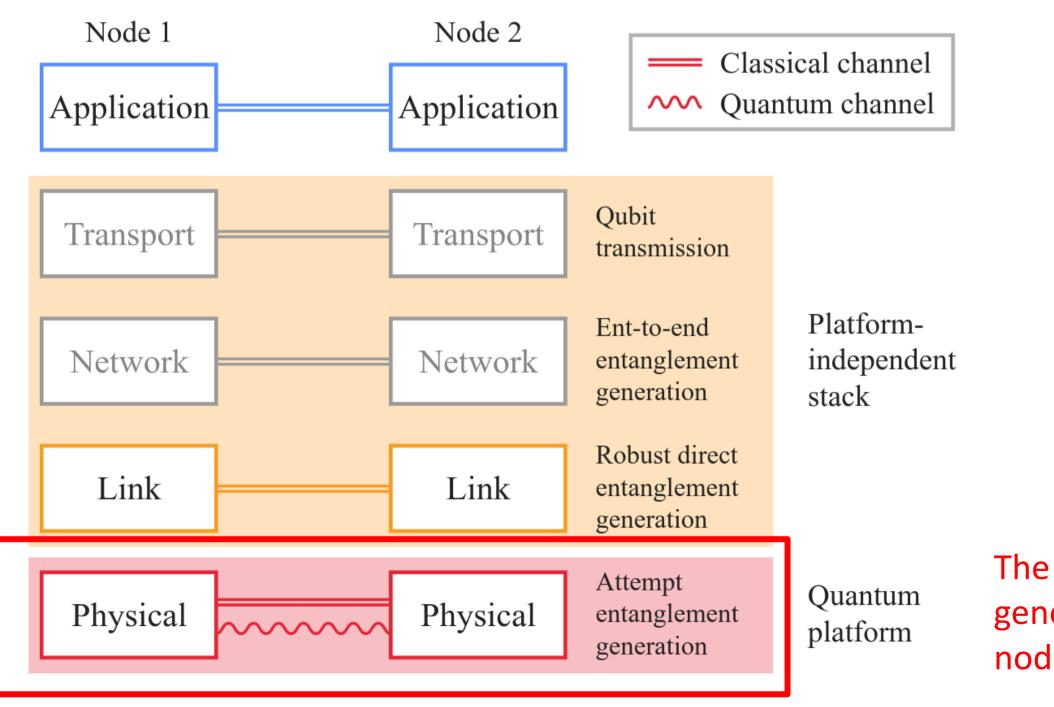
1

# Protocols for scheduling the generation and distribution of quantum entanglements in a quantum network



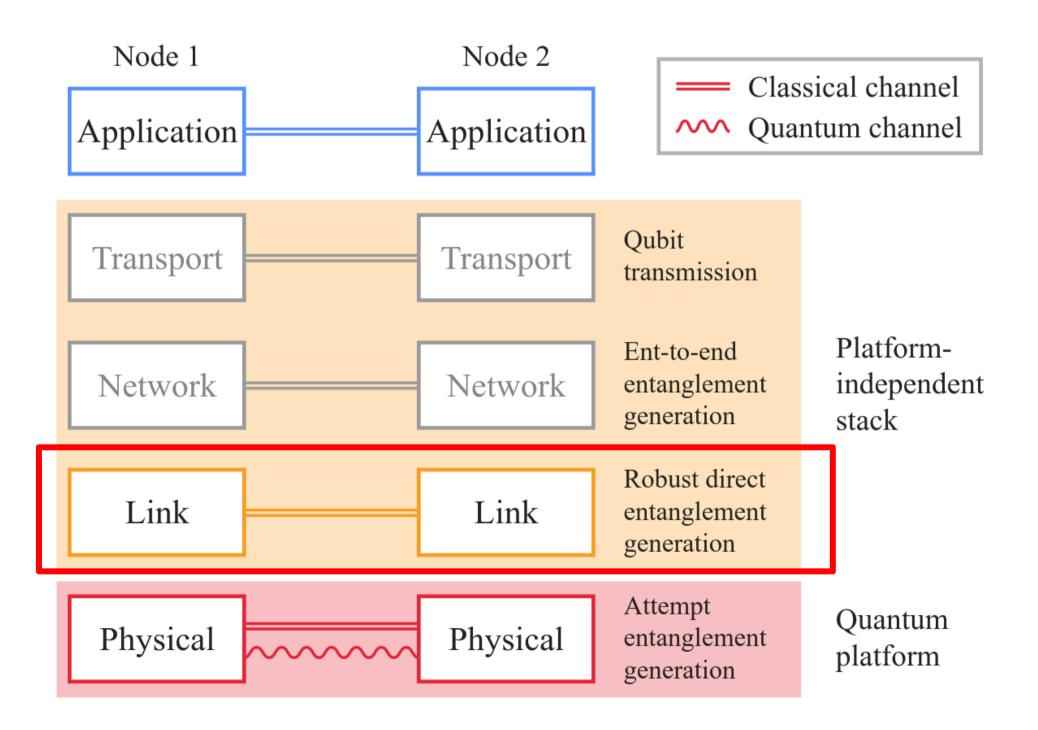
Pompili, Matteo, et al. "Experimental demonstration of entanglement delivery using a quantum network stack." npj Quantum Information 8.1 (2022): 121.

# Protocols for scheduling the generation and distribution of quantum entanglements in a quantum network



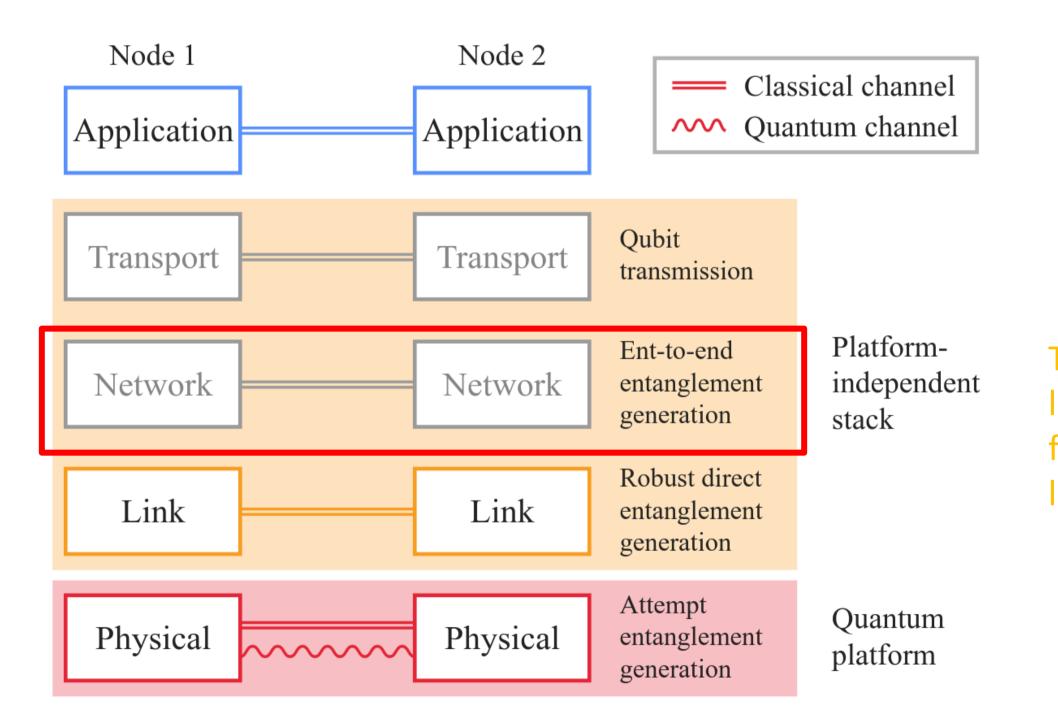
The **Physical Layer** attempts to generate entanglements between two nodes in a well-defined time slot.

# Protocols for scheduling the generation and distribution of quantum entanglements in a quantum network



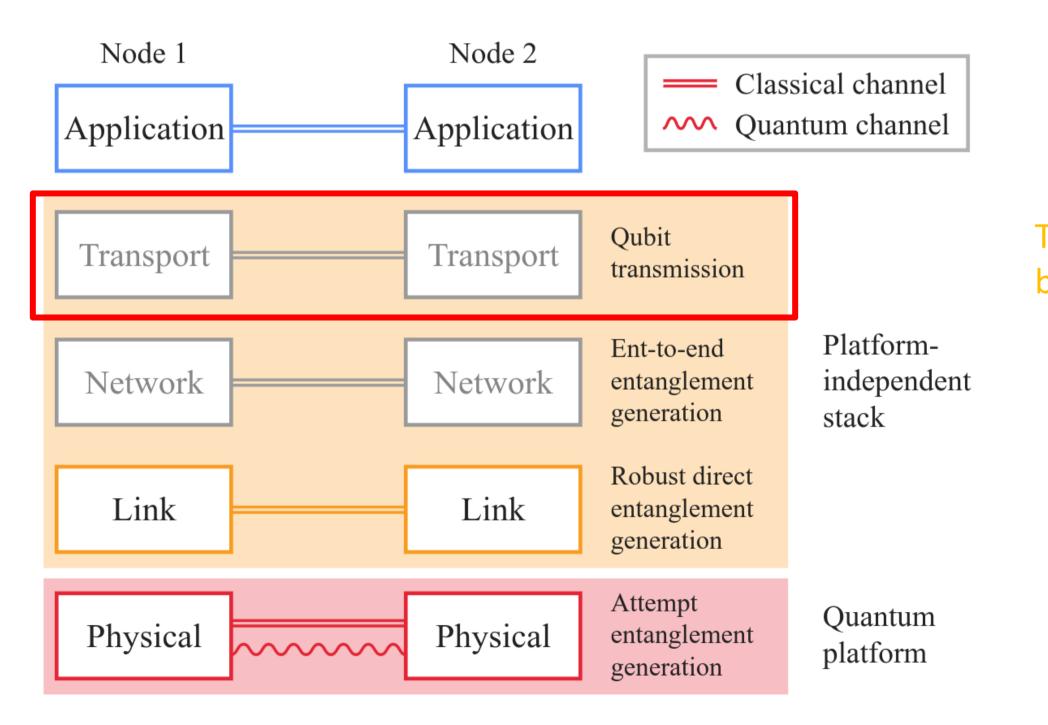
- The **Link Layer** manages to generate robust entanglements:
- Receives generation requests
- Perform fidelity evaluation
- Scheduling generation

# Protocols for scheduling the generation and distribution of quantum entanglements in a quantum network



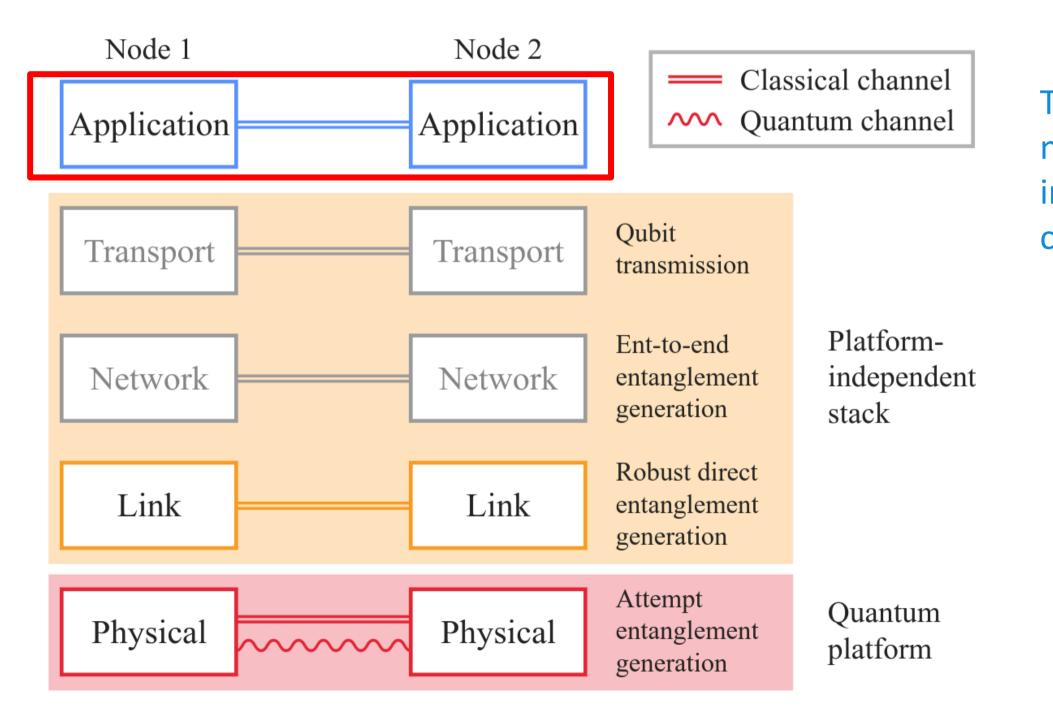
The **Network Layer** is for producing long-distance entanglements using functionalities provided by the link layer.

# Protocols for scheduling the generation and distribution of quantum entanglements in a quantum network



The **Transport Layer** transmits qubits by using the teleportation process.

# Protocols for scheduling the generation and distribution of quantum entanglements in a quantum network



The **Application Layer** controls the network abstractly to distribute quantum information and to perform joint quantum computation.

### **Future Directions**





### What challenge do you want to tackle in the NISQ era?





Quantum Internet

